

NSWI090: Computer Networks

<http://www.ksi.mff.cuni.cz/~svoboda/courses/202-NSWI090/>

Lecture 9

Addresses and Addressing I

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Lecture Outline

Addresses and addressing

- L2: **MAC** addresses
- L3: **IP** addresses
- L4: **port** numbers
- L7: **URI** identifiers

Addressing at L2

MAC addresses (**hardware addresses**)

- Used at L2
 - More precisely at the **MAC sublayer (Media Access Control)**
- Assigned to **network interface controllers**
 - I.e., individual L2 interfaces of end nodes as well as routers
 - Bridges / switches only have one MAC address, if any
 - In case they should explicitly be accessible as ordinary devices
 - Not because of fulfilling their L2 tasks (filtering / forwarding)
- Types of addresses
 - Standard **unicast**, but also **multicast** and **broadcast**
- **Different technologies** may have different mechanisms
 - However, **sharing of address spaces** may be desirable
 - So that **Wi-Fi and Ethernet** can coexist in one network

Addressing at L2

MAC addresses (cont'd)

- Must be **locally unique** within a given network
 - So that senders can **identify** the intended recipients
 - And these recipients are able to **recognize** their frames
 - Since everything is / may be delivered to everyone
 - Because all nodes are mutually visible and reachable
- **Assignment strategies**
 - **Locally** administered addresses would suffice
 - However, **globally unique addresses** simplify everything
 - Otherwise newly connected devices would need to be treated
 - So that they do not potentially use conflicting addresses
- ⇒ **globally unique burned-in addresses assigned by device manufacturers** are primarily used in practice
 - Exceptions exist, though

EUI Addresses

EUI numbering systems (Extended Unique Identifier)

- **EUI-48** (48 bits, 6 bytes, \approx 281 trillion addresses)
 - Formerly denoted as **MAC-48**
 - Which is inappropriate since MAC denotes the entire sublayer
 - And so this label should no longer be used since it is obsolete
 - **Notation**
 - Six hexadecimal numbers separated by hyphens or colons
 - E.g.: FC-77-74-19-41-1E or FC:77:74:19:41:1E
 - Deployment: **Ethernet**, **Wi-Fi**, **Bluetooth**, **ATM**, ...
- **EUI-64** (64 bits, 8 bytes)
 - Newer version with larger address space
 - **Conversion** of EUI-48 possible by adding FF-FE in the middle
 - E.g.: FC-77-74-FF-FE-19-41-1E
 - Deployment: **FireWire** (IEEE 1394), ...

EUI Addresses

Internal structure

- **Organizationally Unique Identifier (OUI)**
 - **Higher 3 bytes** (24 bits)
 - Describes a particular **vendor or manufacturer**
 - E.g.: FC-77-74 (Intel), 90-F6-52 (TP-Link), ...
- **Interface number**
 - Lower 3 bytes (24 bits) in case of EUI-48
 - Lower 5 bytes (40 bits) in case of EUI-64
 - **Serial number** of a given **network interface controller**

OUI Identifiers

OUI component details

- **M bit (I/G bit)** = the least significant bit of the first byte
 - **0** = **individual address** (unicast)
 - **1** = **group address** (multicast, broadcast)
- **X bit (U/L bit)** = the second least significant bit of the first byte
 - **0** = **Universally Administered Addresses (UAA)**
 - Intended for globally unique addresses
 - **1** = **Locally Administered Addresses (LAA)**
- Examples
 - **Individual (M = 0) and universal (X = 0)**
 - FC-77-74-19-41-1E: in binary 11111100-...
 - **Group (M = 1) and local (X = 1)**
 - FF-FF-FF-FF-FF-FF (L2 broadcast): in binary 11111111-...

OUI Identifiers

Observations

- Why M and X bits are the **two least significant bits**?
 - **Ethernet** uses **Big Endian** for the individual **bytes**
 - But **Little Endian** for individual **bits** within these bytes
 - Therefore the first two transmitted bits become M and X
 - And so various address modes can be distinguished easily
- Both M and X bits are **set to zeros** in **OUI identifiers** as such
 - Only when they are exploited in EUI addresses...
 - ... they can then be changed as required and so their special meaning activated and utilized

OUI Identifiers

Organizationally Unique Identifier (OUI) (cont'd)

- Unique **vendor, manufacturer, or other organization** identifier
 - **Purchased** from the **IEEE Registration Authority**
 - One organization may actually have multiple OUIs at a time
- Various usages
 - Our **EUI-48 and EUI-64 addresses**
 - But also **SNAP Protocol Identifiers** (IEEE 802.2 LLC extension)
 - ...
- **Online list:** <http://standards-oui.ieee.org/oui/oui.txt>
 - Current status (April 2021)
 - Almost 30 thousand OUIs are assigned
 - Huawei \approx 1100, Cisco \approx 1060, Apple \approx 890, Samsung \approx 690, ...
 - TP-Link \approx 160, Technicolor \approx 80, D-Link \approx 60, Zyxel \approx 40, ...
 - More than 17 thousand organizations only have a single OUI

Addressing at L4

Port numbers

- Assigned to **access points** between L4 and L7 in order to allow **end-to-end** communication of individual **application entities**
 - Within a given end node
 - **Separately for each transport protocol (TCP, SCTP, DCCP, UDP)**
 - Yet in practice usually the same for all these protocols
- Necessary for both **incoming and outgoing** directions
- Allow for the identification of **transport connections**
 - Tuple (sender $IP_1:port_1$, protocol, recipient $IP_2:port_2$)
 - **Target** of the **outgoing** transmission ($IP_2:port_2$, protocol)
 - **Source** of the **incoming** transmission ($IP_1:port_1$, protocol)
- **Requirements**
 - Ports must be **unique, abstract, implicit and static**

Ports

Port numbers

- **16-bit** long integer numbers: permitted **values 0 – 65535**
- System and registered ports maintained by **IANA**
 - **Internet Assigned Numbers Authority**
- Online table of **port numbers** and **service names**
 - <https://www.iana.org/assignments/service-names-port-numbers/>

Types of ports

- **System Ports (Well Known Ports) (0 – 1023)**
 - **Assigned to one purpose** (usually system-oriented services)
 - Should not be used for any other purpose
 - Examples
 - FTP (21), SSH (22), SMTP (25), DNS (53), HTTP (80), POP3 (110), NTP (123), IMAP (143), HTTPS (443), ...

Ports

Types of ports (cont'd)

- **User Ports (Registered Ports)** (1024 – 49151)
 - **Assigned to one purpose** again
 - But may freely be used for any other purpose
 - Examples
 - Registered: MySQL (3306), PostgreSQL (5432), Redis (6379), Neo4j (7474), MongoDB (27017), ...
 - Not registered: Riak (unsigned 8087 and 8098), Cassandra (assigned 7000), ...
- **Dynamic Ports (Private Ports)** (49152 – 65535)
 - Not assigned, available for unrestricted usage
 - Usually for **outgoing transmissions**

Addressing at L7

Requirements and expectations

- **Various kinds of objects** need to be identified at L7
 - Web pages, files, e-mail addresses, publications, ...
- Two aspects actually need to be covered
 - **Identification**
 - So that objects of one kind can mutually be distinguished
 - At least **locally** (within a given end node) but also **globally**
 - **Location**
 - In terms of a particular node where such objects can be found
 - It make sense to logically decouple both these aspects
- **Each application** may have its own proprietary **naming system**
 - Yet it makes sense to pursue **unification** and **coordination**
 - As well as to recycle approaches that already exist

URI Framework

Uniform Resource Identifier (URI)

- Generic, federated and extensible **naming system**
 - Allows to identify basically anything
 - Including real-world objects (people, places, concepts, ...)
- Types of identifiers
 - **Uniform Resource Locator (URL)**
 - **Web resource reference** (web address)
 - Specifies particular location as well as retrieval mechanism
 - **Uniform Resource Name (URN)**
 - **Globally unique persistent resource identifier**
 - Does not imply any location, not widely used
 - **Uniform Resource Characteristics (URC)**
 - Description of **meta data** about URLs or URNs (citations, ...)
 - Never standardized, not even implemented

URI Examples

Sample URLs

- **http** scheme (**H**ypertext **T**ransfer **P**rotocol)
 - Addresses of web pages or other resources
 - E.g.: `http://www.mff.cuni.cz/en/index.php?page=people#5460`
- **ftp** scheme (**F**ile **T**ransfer **P**rotocol)
 - Paths to files or directories accessible using the FTP protocol
 - E.g.: `ftp://svoboda:password@ulita.ms.mff.cuni.cz/`
- **file** scheme
 - Host-specific paths on local or remote **file systems**
 - E.g.: `file:///home/svoboda/NSWI090/Lecture-03-Layers.pdf`
- **mailto** scheme
 - **E-mail addresses** including additional parameters
 - E.g.: `mailto:svoboda@ksi.mff.cuni.cz?subject=NSWI090`

URI Examples

Sample URLs

- **tel** scheme
 - **Telephone numbers**
 - E.g.: `tel:+420-951-554-250`
- **sip** scheme (**Session Initiation Protocol**)
 - Participants of multimedia sessions such as **voice calls** (VoIP, ...)
 - E.g.: `sip:martin.svoboda@mff.cuni.cz`
- **jdbc** scheme (**Java Database Connectivity**)
 - Connections to **relational databases** from Java applications
 - E.g.: `jdbc:postgresql://nosql.ms.mff.cuni.cz:5432/database`

URI Examples

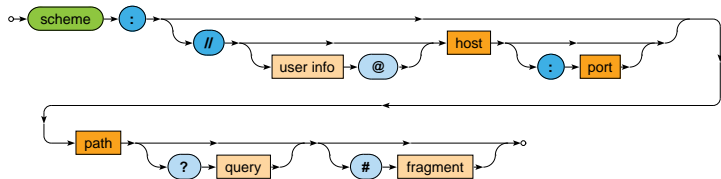
Sample URNs

- **isbn** namespace (**International Standard Book Number**)
 - Printed or electronic **books**
 - E.g.: `urn:isbn:9780132126958`
- **issn** namespace (**International Standard Serial Number**)
 - Printed or electronic **serial publications** (journals, ...)
 - E.g.: `urn:issn:0302-9743`
- **ietf** namespace (**Internet Engineering Task Force**)
 - IETF family of RFC, STD, FYI, and BCP documents
 - E.g.: `urn:ietf:rfc:2648`
- **iso** (**International Organization for Standardization**)
 - **Standards** and other technical specifications
 - E.g.: `urn:iso:std:iso-iec:9075:-1:ed-5:en`

URI Structure

Generic syntax

- May further be restricted by particular schemes
- **Syntax diagram** (for context-free grammars)



- **Green boxes: literals**
 - Expected to be replaced with a particular value
- **Blue boxes: constants** (preserved as they are)
- **Orange boxes: non-terminals**
 - Unfolded to more complicated fragments

URI Schemes

URI components

- **Scheme**: case-insensitive **scheme name** (usually lower-case)
- **Authority**
 - **User info**: authentication tokens such as **name** and **password**
 - Deprecated for security reasons
 - **Host**: **domain name**, **IP address**, or a different registered name
 - **Port**: transport layer **port**
- **Path**: usually **hierarchical path** with individual segments
- **Query**: usually parameters in a form of **attribute / value pairs**
- **Fragment**
 - Reference to a **secondary resource** related to the primary one
 - E.g.: anchors in HTML pages, classes in OWL ontologies, ...

URI Framework

Both **schemes** and **namespaces** are registered with **IANA**

- **URI schemes**

- <https://www.iana.org/assignments/uri-schemes/>
- Current status (April 2021)
 - \approx 100 permanent, 230 provisional, and 10 historical schemes

- **URN namespaces**

- <https://www.iana.org/assignments/urn-namespaces/>
- Current status (April 2021)
 - \approx 70 namespaces

Internationalized Resource Identifier (IRI)

- Just an extended version of the traditional URIs
- Allows to use most of **Unicode characters**
 - And so Chinese, Japanese, Korean or other national characters

Addressing at L3

IP addresses

- Primary objective is **addressing of nodes** as a whole
- In spite of that...
 - IP addresses are actually assigned to their **network interfaces**
 - I.e., individual L3 interfaces of end nodes as well as routers
 - Note that end nodes used to usually have only one IP address
 - But nowadays, **multi-homed** hosts with more IPs are common
- Another observation...
 - **Routing algorithms perceive networks as a whole**
 - ⇒ **networks themselves** must also be **uniquely identified**
 - Though two separate network / node identifiers could work
 - **Internally structured atomic IP addresses** are more practical
 - Otherwise hop-to-hop **routing** and **forwarding** would not work

IP Addresses

Requirements on IP addresses

- **Globally unique** within the whole system of networks
- And internally structured...

Internal components

- **Network part (Network ID)**
 - **Prefix** of the whole address
 - I.e., certain part from the beginning
 - **Uniquely identifies a given network** as a whole
 - Determines affiliation of nodes to a particular network
- **Relative part (Host ID)**
 - Remaining part of the whole address
 - **Uniquely identifies a given node** within a given network

Assignment Principles

Observations = rules that must be followed

- **Two nodes in the same network...**
 - Must have the same network parts
 - And different relative parts
 - So that they can be mutually distinguished
- **Two nodes in different networks...**
 - Must have different network parts
 - So that we can detect they belong to different networks
 - And as for relative parts, they do not matter
 - They may be the same as well as not

Assignment Principles

Block principle

- (1) **network as a whole** must first be assigned with a whole contiguous **block of addresses**
 - I.e., set of IP addresses belonging to a specific range
 - All with the same prefix (network part)
 - **Assignment process** must be **globally coordinated**
 - **IANA** and regional providers
- (2) only than **individual nodes** can be given their addresses
 - Of course, from this range
 - Manually or using **DHCP** or similar protocols

Consequence

- **Unused addresses cannot be used by anyone else!**
 - This may lead / actually led to unacceptable wasting

IPv4 Addresses

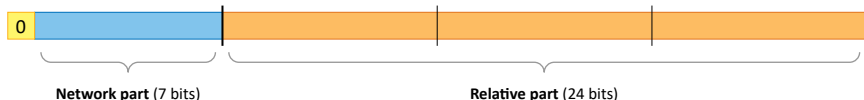
IPv4 addresses

- **4 bytes** (32 bits) – potentially \approx 4 billion values
 - Which is certainly not much from today's perspective
 - But was initially thought of as generously sufficient
- **Notation**
 - Four decimal numbers, one for each byte, separated by dots
 - E.g.: 195.113.19.170
- **Internal structure: network and relative parts**
 - Where the divide between the parts should be located?
 - 3 possible **divide placements** are possible
 - This gives us **Classes A, B and C** of IP addresses
 - If more options exist, though, how to recognize them?
 - Since it must be possible just from the IP address itself

Classes of IPv4 Addresses

Class A

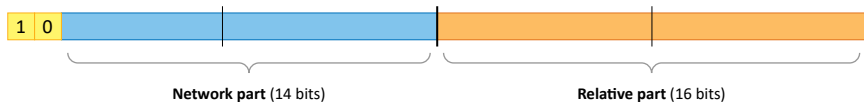
- Divide is positioned **after the first byte**
 - \approx 128 networks
 - Each with \approx 17 million addresses
- Overall **range**: **0.0.0.0** – **127.255.255.255**
 - Covers 1/2 of the entire space
- **Highest bit** is 0
- Suitable for **very large networks**



Classes of IPv4 Addresses

Class B

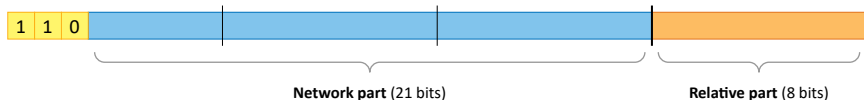
- Divide is positioned **after the second byte**
 - \approx 16 thousand networks
 - Each with \approx 66 thousand addresses
- Overall **range**: **128.0.0.0** – **191.255.255.255**
 - Covers 1/4 of the entire space
- **Highest bits** are **10**
- Suitable for **medium-sized networks**



Classes of IPv4 Addresses

Class C

- Divide is positioned **after the third byte**
 - \approx 2 million networks
 - Each with \approx 256 addresses only
- Overall **range**: **192.0.0.0** – **223.255.255.255**
 - Covers 1/8 of the entire space
- **Highest bits** are **110**
- Suitable for **very small networks**



Classes of IPv4 Addresses

Apparently the **available space** is not yet fully utilized...

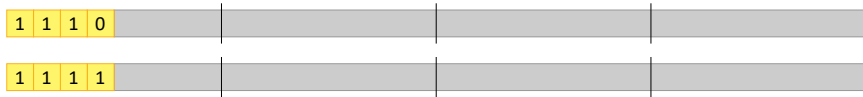
- The rest is covered by **Classes D and E**
 - They have **specific usage** and they are **internally unstructured**

Class D: multicast addresses

- 224.0.0.0 – 239.255.255.255 (1/16 of the entire space)
- Highest bits are **1110**

Class E: reserved for **future extensions**, but never used

- 240.0.0.0 – 255.255.255.255 (1/16 of the entire space)
- Highest bits are **1111**



Special Addresses

Special addresses

- Certain addresses (even whole ranges) have **dedicated usage**
 - Basic principle: bit 0 = **this**, bit 1 = **all**

Nodes

- Self node: **0.0.0.0**
 - Useful when standard unicast address is not yet known
- Node in a local network: **0.X.Y.Z**, **0.0.X.Y**, or **0.0.0.X**

Networks

- Network as a whole: **X.0.0.0**, **X.Y.0.0**, or **X.Y.Z.0**

Broadcasts

- **Targeted** broadcast: **X.255.255.255**, **X.Y.255.255**, or **X.Y.Z.255**
- **Limited** broadcast: **255.255.255.255**

Special Addresses

Loopback

- 1 A-block: **127.X.Y.Z**
 - **127.0.0.1** is in particular usually used
 - But the whole block is in fact available

Link local addresses

- 1 B-block: **169.254.X.Y**
 - Auto-configuration when standard address cannot be obtained

Private addresses

- 1 A-block: **10.X.Y.Z**
- 16 B-blocks: **172.16.X.Y** – **172.31.X.Y**
- 256 C-blocks: **192.168.0.X** – **192.168.255.X**

Multicast Addresses

Multicast transmissions

- Intended recipients are all nodes in a given group
 - This group is predefined or created dynamically on demand
 - **IGMP (Internet Group Management Protocol)**
- Certain blocks are assigned and approved by **IANA**
 - <https://www.iana.org/assignments/multicast-addresses/>
 - Current status (April 2021)
 - \approx 60 static multicast addresses (beside other)
- **Class D addresses** are used
 - Start with **1110** and they are **internally unstructured**
 - I.e., there is no network / relative part



Multicast Addresses

Static groups (well known groups)

- Node membership is fixed and given in advance
- Range: **224.0.0.0** – **224.0.0.255** (\approx 256 addresses)
 - Reserved for **routing and other low-level protocols**
 - Topology discovery, maintenance, ...
 - **224.0.0.1**: all hosts in a given network
 - **224.0.0.2**: all routers in a given network
 - ...

Dynamic groups

- **Global**: **224.0.1.0** – **238.255.255.255** (\approx 252 million)
 - Group scope can span multiple different networks
- **Local**: **239.0.0.0** – **239.255.255.255** (\approx 17 million)
 - Group scope is limited to a particular network

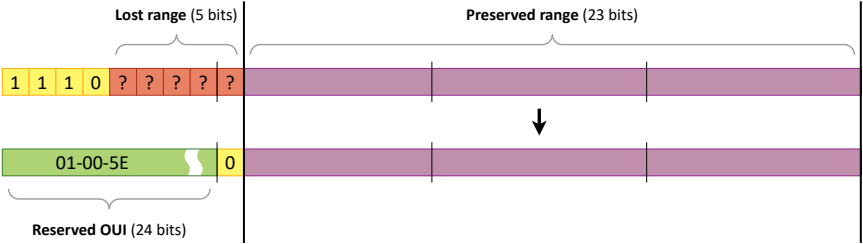
Multicast Addresses

Translation of multicast addresses (IP \rightarrow EUI-48)

- L3 multicast **IP address**
 - 4 bits at the beginning are fixed \Rightarrow only **28 bits are relevant**
 - Unfortunately, **only the last 23 bits** can be considered
 - And so the remaining 5 bits in the middle must be truncated
 - \Rightarrow 32 different addresses are mapped to the same value!
- L2 multicast **EUI-48 address**
 - 00-00-5E reserved **OUI** is used (24 bits)
 - Of course, altered to 01-00-5E so that M bit = 1
 - The 25th bit is set to 0
 - The remaining 23 bits are taken from the original IP address

Multicast Addresses

Translation of multicast addresses (IP → EUI-48) (cont'd)



Allocation Strategies

Basic principle

- **Whole blocks** of addresses are / must be assigned to networks
- These blocks correspond to our **classes**
 - One **Class A** block \approx 17 million individual addresses
 - One **Class B** block \approx 66 thousand individual addresses
 - One **Class C** block \approx 256 individual addresses

Strategies

- Formerly: **one closest larger block principle**
 - E.g.: 1 Class B address for 1000 requested individual addresses
 - Extremely wasteful approach...
- Later: **multiple closest smaller blocks principle**
 - E.g.: just 4 Class C addresses for the same requested number
 - Better, but still not enough...

Lack of IPv4 Addresses

Temporary **mitigating solutions**

- 1985: **Subnetting**
 - One larger network is divided into separate sub-networks
- 1988: Allocation mechanism
 - **One larger block** → **more smaller blocks** principle
- 1993: **CIDR (Classless Inter-Domain Routing)**
 - Original concept of IP address classes is entirely dropped
- 1994: **Private addresses**
 - Usage of private IPv4 addresses instead of globally unique ones
 - Requires **NAT (Network Address Translation)**

Permanent solution

- 1995: IPv6 protocol and its **IPv6 addresses**
 - **6 bytes** instead of 4 bytes ⇒ significantly **larger address space**

Lecture Conclusion

L2: **MAC** addresses

- **EUI-48** and EUI-64 numbering systems, **OUI**, M and X bits

L3: **IP** addresses

- **Network and relative** parts
- **Classes A, B, C, D, and E**
- Special and **multicast** addresses

L4: **port** numbers

- System / User / Dynamic ports

L7: **URI** identifiers

- URI (IRI) framework: **URL** locators, **URN** names