NSWI090: Computer Networks

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Lecture 1

Basic Concepts and Paradigms

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Lecture Outline

Introduction and motivation

- Basic concepts
- Network paradigms
- Telecommunication / computer networks

Motivation

Ultimate objective

- Transmission of data
 - I.e., not processing
 - Algorithms, data structures, time complexities, ...
 - As well as not storing
 - Database systems, logical models, query languages, ...

Major requirement

- Transmitted data should not change
 - However...

Relevant questions

• Distance, volume, velocity, latency, reliability, ...

Transmission path

- Path between the sender and recipient in a network
- Includes any means of transmission
 - Regardless of technical implementation
 - Wired (metallic, optical), wireless
 - At different levels of abstraction
 - Physical, virtual

Transmission channel

One-way transmission path in telecommunication networks

Transmission circuit

- Two-way transmission path
 - Physical, emulated / virtual

Simplex mode

- Allows for transmissions in one direction only
 - E.g., television and radio broadcasting

Half-duplex mode

- Allows for transmissions in **both directions**, but **not simultaneously**
 - Communicating entities need to take turns

Full-duplex mode

• Allows for transmissions in **both directions** at the same time

Unicast transmissions

- One sender, one intended recipient
 - The most common situation

Anycast transmissions

Any node in a given predefined group

Multicast transmissions

Every node in a static / dynamically created group

Broadcast transmissions

• Any reachable node (in a given network, ...)

Telecommunication / computer network / infrastructure

• System of nodes mutually interconnected by transmission paths formed by network elements

Nodes

- Anything wanting to communicate
 - Servers, terminals, personal computers, laptops, ...
 - Printers, mobile phones, smart TVs, ...

Network elements

- Active
 - Repeaters, switches, routers, gates, ...
 - Switchboards, controllers, transceivers, ...
- Passive

Networks provide services to their users

- Television and radio broadcasting
- Voice and video calls
- Video on demand
- Electronic mail
- Messaging platforms
- File transfers

...

Web pages browsing

Services (cont'd)

Provided at different levels of abstraction

- Internal implementation details are concealed
- Services can be built on top of other services
 - HTTP web page request at L7 \rightarrow reliable connection-oriented TCP at L4 \rightarrow best effort IP at L3 \rightarrow Ethernet at L2

Different requirements and expectations

Different principles, protocols, and technologies are needed

Network Paradigms

Fundamental questions to be figured out

- Stream vs. block transmissions
 - In what form will the data be transmitted?
- Circuit vs. packet switching
 - How will the data reach the intended recipient?
- Connection-oriented vs. connectionless transmissions
 - Do we need to agree on the communication in advance?
- Reliable vs. unreliable transmissions
 - What level of transmission reliability is required?
- Guaranteed vs. non-guaranteed services
 - Will sufficient resources be available all the time?

Data Transmission

In what form will the data be transmitted?

Stream transmissions

Data is sent as a continuous stream without any logical division

Block transmissions

• Data is divided into reasonably large blocks, these blocks are then sent one by one

Stream Transmissions

Streams

- Data is transferred as a continuous unstructured stream
 - At the level of individual symbols (bits, bytes, words, ...)



- Features
 - Preserves order (works as a FIFO queue)
 - Suitable for point-to-point connections
 - I.e., in case there is only one intended recipient
- Examples
 - L1: Ethernet, Wi-Fi, ...
 - L4: TCP (virtual streams for L7 over L3 blocks)

Block Transmissions

Blocks

- Data is divided into reasonably large units called blocks
 - Fixed size
 - Variable size
 - Restrictions on maximal / minimal lengths may be in effect, too



- Features
 - Blocks may not directly follow each other
 - Nothing is remaining to be sent at a moment
 - Intentional delays between the blocks may be needed
 - Order may not be preserved

Block Transmissions

May (but also may not) be enriched with additional metadata

- Sender / recipient addresses
 - So that blocks can be correctly forwarded when a given transmission path is not direct

Transmission identification

 So that multiple different transmissions can be mutually distinguished on a shared path

Block ordinal numbers

• ...

So that the original data can once again be reconstructed

Block Transmissions

Main tasks

- Internal structure
 - Header, body, footer
- Combinations of both the approaches
 - Blocks over streams (framing)
 - Blocks over blocks (encapsulation)
 - Streams over blocks

Names of blocks

- Depend on particular layers / technologies / protocols
 - L2: Ethernet frames, ATM cells, ...
 - L3: IPv4 datagrams, IPv6 packets, ...
 - L4: TCP segments, UDP datagrams, ...
 - L7: HTTP messages (requests, responses), ...

Switching Mechanisms

How will the data reach the intended recipient?

Circuit switching

- Transmission circuit is created first, all data is then sent through this circuit
- Based on reserved capacity
 - Mainly in telecommunication networks

Packet switching

- Each block of data is routed and forwarded separately
- Based on shared capacity
 - Mainly in computer networks
 - Very revolutionary idea in its time (in 1960s)
 - Aims at robustness and resilience against failures and outages

Circuit Switching

Basic mechanism

Transmission circuit is created on demand

- I.e., transmission path is found and laid out
 - Physically interconnected
 - Marked virtually at higher layers
- All data is then sent directly through this circuit

Examples

- PSTN (traditional Public Switched Telephone Network)
 - Originally manual switching by operators → later on automated electronic switching → fully digital switching nowadays

Circuit Switching

Features and consequences

- Illusion of a direct connection
 - Though a circuit will most likely in reality be laid out across multiple network elements
- \Rightarrow low and constant latency
- \Rightarrow preserves order
- Effect of reserved capacity
 - No one else is entitled to use a given circuit
- \Rightarrow enables **guarranteed** transmissions
- Supports both streams and blocks
- Always is connection-oriented

Switching node logic

- Each interface (path) has its own inbound / outbound buffer
- Whenever a new block is received
 - It is stored within a given inbound buffer
- Waiting blocks are perpetually iterated over and processed
 - Routing decision is made for a given block
 - It is then moved to the corresponding outbound buffer
 - Finally, it is send as soon as it is possible

Examples

- L2: Ethernet switches, ...
- L3: IP routers, ...

Only possible for block transmissions

- \Rightarrow the following pieces of metadata is necessary
 - Recipient identification
 - So that the intended recipient can be located as well as the individual ongoing transmissions mutually distinguished
 - Sender identification
 - So that error messages can be sent back in case of failures

Available capacity is shared and limited

- Transmission capacity of paths
- Computing (switching) capacity of switching nodes
 - Buffer sizes, processor throughput
- \Rightarrow significantly higher, variable, and unpredictable latency
 - Which also depends on the current load
 - I.e., it can never be estimated in advance
- \Rightarrow available capacity may be insufficient
 - Excessive blocks will then need to be discarded

Connection Schemes

Do we need to agree on the communication in advance?

Connection-oriented transmissions

- Both the communicating parties need to establish, maintain, and eventually terminate a connection with each other
- Order of data is / must be preserved
- Unique connection ID is usually assigned

Connectionless transmissions

- Communication is based on sending separate messages
- Order of data is not preserved

Connection-Oriented Transmissions

Connection-oriented approach

- Connection is established
 - This means that...
 - Both the parties really exist
 - They are able to locate each other
 - They **agree** with the communication
 - Communication parameters can be negotiated
 - Transmission path can be laid out, resources allocated, ...
- The actual data is then transferred
- Connection is eventually terminated
 - E.g., resources are returned, marked circuits are released, ...

Connection-Oriented Transmissions

Analogy

- Traditional voice call over a telephone network
 - Phone number is dialed \to call is established \to parties are talking to each other \to call is hanged up and terminated

Examples

- L2: ATM (Asynchronous Transfer Mode)
- L4: TCP (over a connectionless IP at L3)
 - Ordering of segments needs to be enforced
- L7: HTTP, SMTP, POP3, ...
 - I.e., anything over TCP from L4

Connection-Oriented Transmissions

Stateful operation

- Communicating parties transition between the states
 - E.g.: closed, established, ...
- Transitions must be correct and coordinated
 - Mutual misunderstandings must be avoided
 - Deadlocks must be avoided
- Non-standard situations must be detected and treated
 - Connection failures, ...

Connectionless Transmissions

Connectionless approach

- Separate messages called datagrams are sent
 - No connection is established
 - Recipient may not exist, may not want to communicate, ...
 - Nothing is established, no termination is required, ...
 - Stateless operation
 - Datagrams are delivered independently on each other
 - They may be routed differently
 - Their order cannot be guaranteed
 - Must contain full recipient identification

Connectionless Transmissions

Analogy

- Traditional postal services
 - Postcards, letters, parcels, ...

Examples

- L4: **UDP**
- L3: IP, ICMP
- L2: Ethernet

Circuit switching is always connection-oriented

Packet switching itself may be ...

- Connection-oriented: virtual circuits
 - Transmission path is laid out only virtually
 - Suitable Virtual Circuit Identifier (VCI) is assigned
 - Individual network elements must note the actual path
 - Blocks are then forwarded based on these VCIs
 - Each is delivered via the same path
 - Example
 - L2: **ATM**
- Connectionless: datagram service

Reliability Levels

What level of transmission reliability is required?

Observations

- Transmissions are never ideal
 - It may always happen that the data will be damaged
 - Entire blocks can be lost
 - Actual data can be corrupted

Reliable transmissions

- Sender considers their duty to take care of the remedy
- Errors are detected and treated

Unreliable transmissions

- Errors are not detected, nor treated
- Transmission simply goes on

Reliable Transmissions

Detection mechanisms

- Parity, checksums, CRC, ...
 - Whatever particular approach is exploited, it is impossible to detect all the possible errors and extents of damage

Remedy options

- Error correction codes
- Repeated transmission
 - Positive acknowledgment when received successfully
 - Negative acknowledgment otherwise
 - I.e., request for a retransmission

Example

• L4: TCP (Go-Back-N or Selective Repeat ARQ)

Reliable Transmissions

Consequences

- Higher number of messages
 - Positive / negative acknowledgments
 - Repeated data transmissions
- Delivery regularity is disrupted
 - Significant delays occur
- Messages are slightly bigger (because of the checksums)
- Higher usage of computing and transmission capacity
 - Sender / recipient nodes
 - The entire mechanism needs to be deployed
 - Network: handling of extra messages
- Reliability is never absolute

Unreliable Transmissions

In other words...

- Reliability is always relative only
 - While insufficient in one case, excessive in another
- It is always connected with a non-zero overhead
- ⇒ unreliable transmissions make sense as well
 - Especially in case of multimedia applications
 - Audio / video, interactive / non-interactive
- **Regularity** of delivery is essential, **low latency** may as well Examples
 - L4: **UDP**
 - L3: IP
 - L2: Ethernet, ATM

Guarantee Options

Will sufficient resources be available all the time?

Guaranteed transmission / service

- Provides such a guarantee for all the currently ongoing data transmissions
 - In terms of computing and transmission capacity
- Resources must be reserved in advance
 - During the connection process
- Realized via circuit switching
- Non-guaranteed transmission / service
 - It may happen that sufficient resources will be missing
 - Realized via packet switching

Guaranteed Transmissions

Exclusive capacity

Solely for the purpose of the given communicating parties

\Rightarrow disadvantages

- Must be high enough to cover the maximum expected load
- Unused capacity cannot be left to anyone else
 - It is forfeited uselessly
- The whole approach is ineffective and expensive
 - Everything must be dimensioned for the sum of maxima

Non-Guaranteed Transmissions

Shared capacity

- Cheaper and more efficient solution
 - Everything can be dimensioned for the average load

\Rightarrow disadvantages

- When sufficient remaining resources are not available
 - Because of...
 - Completely filled buffers
 - Overloaded processor
 - Transmission capacity of individual paths
- Certain packets will need to be discarded!
 - Note that this is the only possible measure
- The question is which...

Non-Guaranteed Transmissions

Best effort principle

- Maximum effort...
 - All packets are delivered as long as it is possible
- ... but uncertain outcome
 - Should packet loss be inevitable, all data is treated equally
 - I.e., there are no rules, no priorities, no criteria
- Examples
 - L3: IP
 - L2: Ethernet, ATM (but also various QoS alternatives)

Quality of service (QoS)

- Anything else compared to the best effort principle
 - The extent of particular guarantees may vary

Quality of Service

Relative QoS – principle of prioritization

- Better conditions are provided for certain kinds of data
 - Based on different priorities

Absolute QoS – principle of reservation

- Guarantees the same conditions regardless of the current situation and load
- Resources must be reserved in advance
 - Reservation request must be rejected when not attainable
- Similar to circuit switching
 - And so has the same advantages, but disadvantages as well

Different Worlds

World of telecommunication networks

- Significantly older
 - Communication was considered a strategic interest
- Smart network, dumb devices paradigm

World of computer networks

Dumb network, smart devices paradigm

Both the worlds traditionally built their own separate networks

- They have always differed greatly in many aspects
- But integration and convergence attempts are intensifying

Telecommunication Networks

Smart network, dumb devices

- All intelligence is concentrated in the network
 - Network elements are often single-purpose
 - Easier and usually central management
 - Expensive, cumbersome, inflexible
- User devices can, therefore, be very simple and foolproof

Preferred transmission characteristics

- Circuit switching
- Connection-oriented
- Reliable
- Guaranteed (QoS)

Telecommunication Networks

Additional observations

- Assumption of insufficient resources
 - There are not enough resources to satisfy everyone
 - Focusing on exclusive resource allocation
 - Afraid to sell unreliable services
- Considered a matter of strategic importance
 - High level of regulation
 - Directive decision-making by national governments
 - Gradually liberalized
 - Monopolies are becoming incumbents
- Network owners and users are different entities

Computer Networks

Dumb network, smart devices

- Intelligence is concentrated in user devices
- Network should be minimalist and as efficient as possible
 - The only goal is to transfer data without understanding it
 - Easily adaptable to changes in user behavior
 - Cheaper, straightforward, flexible

Preferred transmission characteristics

- Packet switching
- Connectionless
- Unreliable
- Non-guaranteed (best effort)

Computer Networks

Additional observations

Availability of resources is not a major limiting factor

- I.e., resources are sufficient
- Technical factors have higher importance than the commercial
- Liberalized from the very beginning
 - Bottom-up approach
 - Standardization and coordination challenges
 - In order to achieve compatibility and interoperability
- Network owners and users are often the same entities

Management of Resources

Moore's law

- Observation that the **number of transistors** in a dense integrated circuit **doubles about every 2 years**
 - Projection of a historical trend
 - In effect for more than 55 years...
 - Originally 1 year, later 1.5 years, now 2 years
- Implication
 - Cost of equivalent computing power drops in half every 2 years

Gilder's law

Transmission capacity triples every 1 year

Disk law

• Storage capacity doubles every 1 year

Lecture Conclusion

Basic concepts

- Stream / block transmissions
- Circuit / packet switching
- Connection-oriented / connectionless transmissions
- Reliable / unreliable transmissions
- Guaranteed / non-guaranteed transmissions
 - Best effort / Quality of Service