DYNAMIC HASHING

NDBI007: Practical Class 4

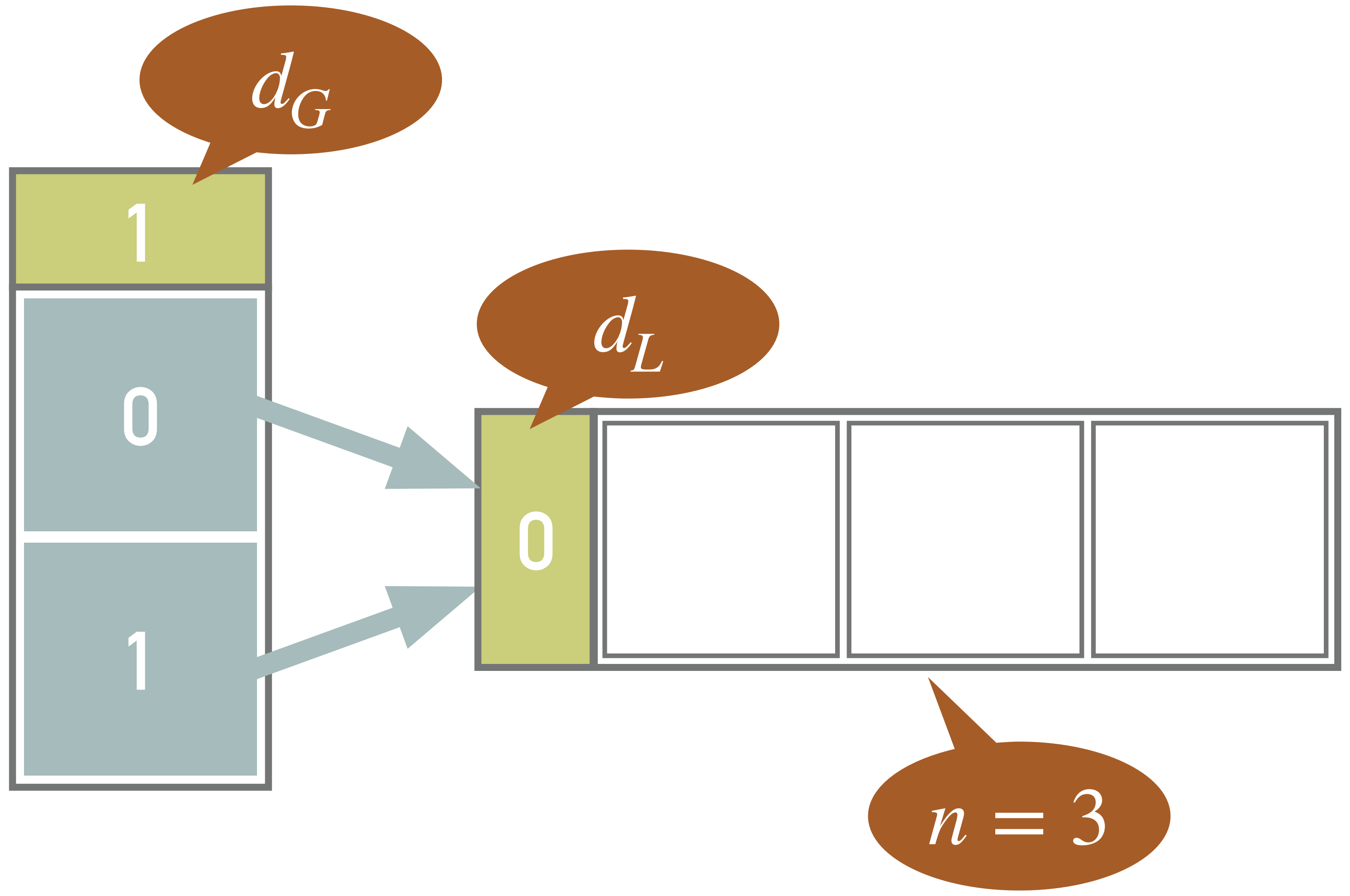
DYNAMIC HASHING

- Static forms of hashing lose its good performance as the table utilisation comes to its maximum
- On the other hand, dynamic hashing algorithms allow to increase the size of the table with increasing number of stored records

- *Fagin*
- *Litwin*
- *LHPE-RL*

FAGIN

- Directory
  - *List of entries* in the *main memory* that points to the pages in the primary file
  - *Global depth* d_G - Number of least significant bits of the hash $h(k)$ needed to *address an entry* in the directory
- Primary file
  - Distributed *collection of pages* stored in the *secondary memory*, i.e., continuous space is not required
  - Each page has a constant size n
  - Each page remembers *local depth* d_L - Number of least significant bits of the hash $h(k)$ *common to all records*
  - $2^{d_G - d_L}$ tells how many directory entries points to the particular page in the primary file



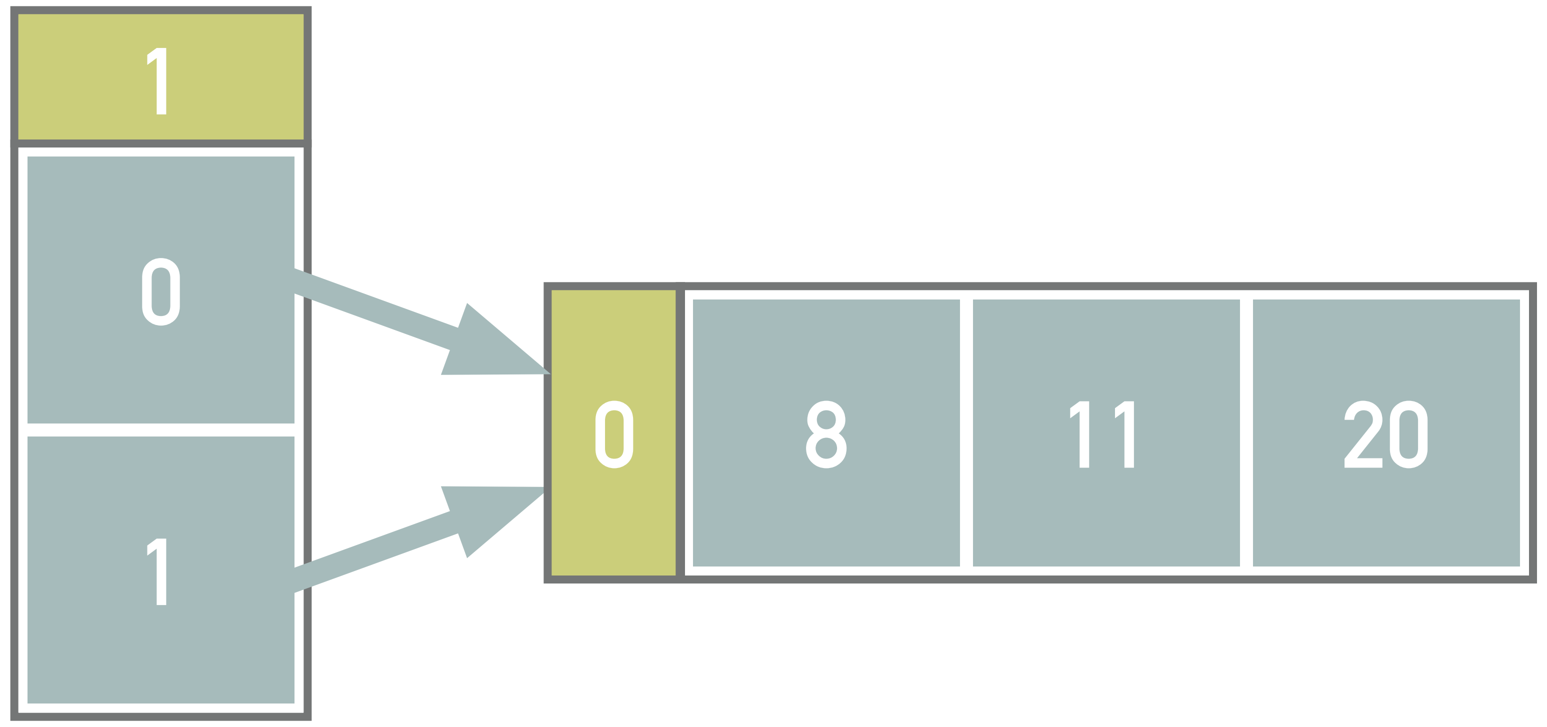
- *Overflowing* causes a *change in the structure* of the directory and primary file
  - $d_L < d_G$ - the particular *page can be split*, i.e., the page is split and d_L incremented
  - $d_L = d_G$ - the *directory* must be *expanded*, i.e., the directory is doubled and d_G incremented

- Inserting/Searching for a record with key k
  - Compute $k' = h(k)$
  - Convert k' into directory entry k'' by leaving the d_G least significant bits
  - The pointer in the corresponding entry points to the page where the record should be inserted/searched

EXAMPLE 1: FAGIN

- Insert records with keys 20, 11 and 8

- $h(20_{10}) = 10100_2$
  - Using the least significant bit of key 20, i.e., 0, the corresponding record is inserted into the page using entry 0

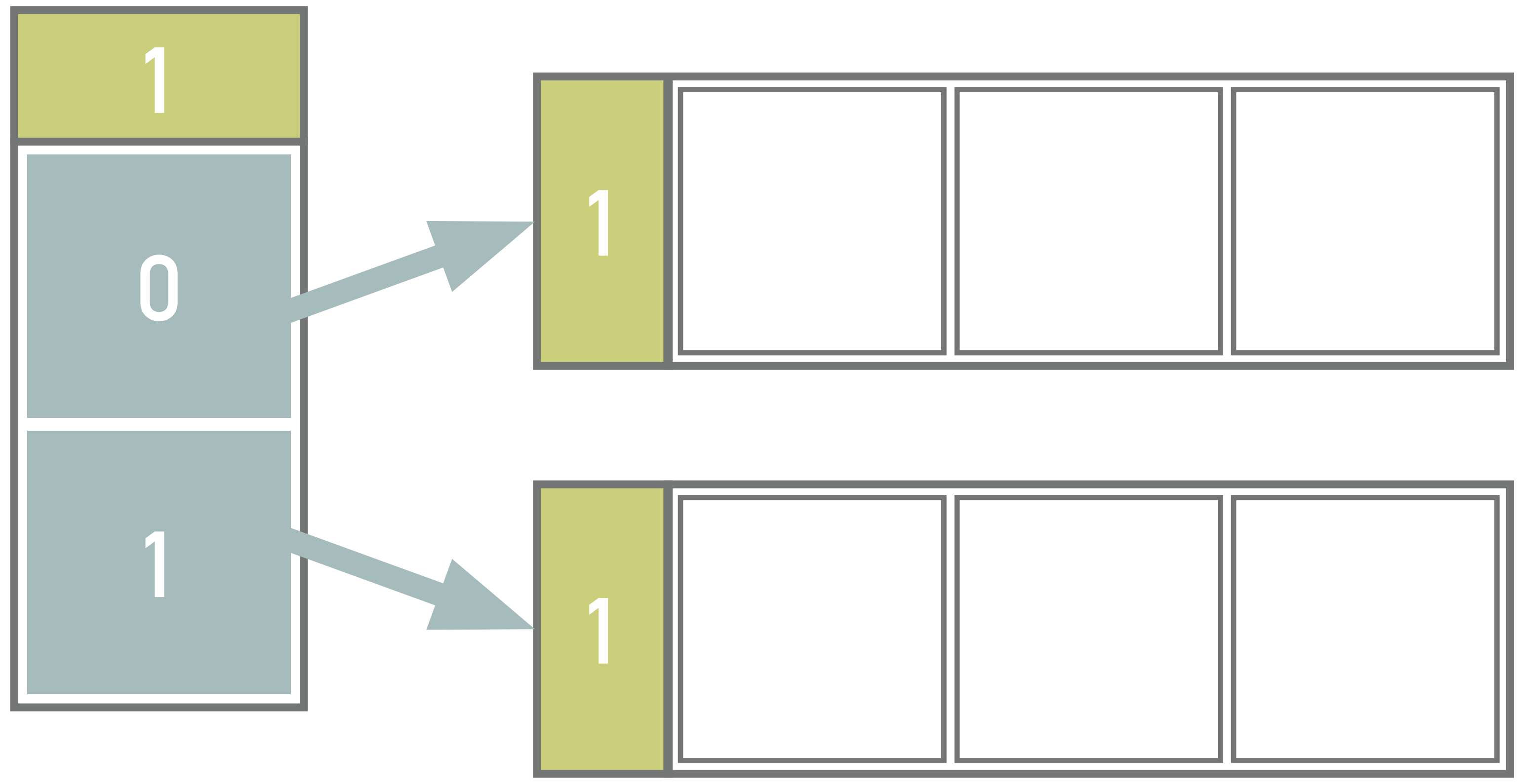


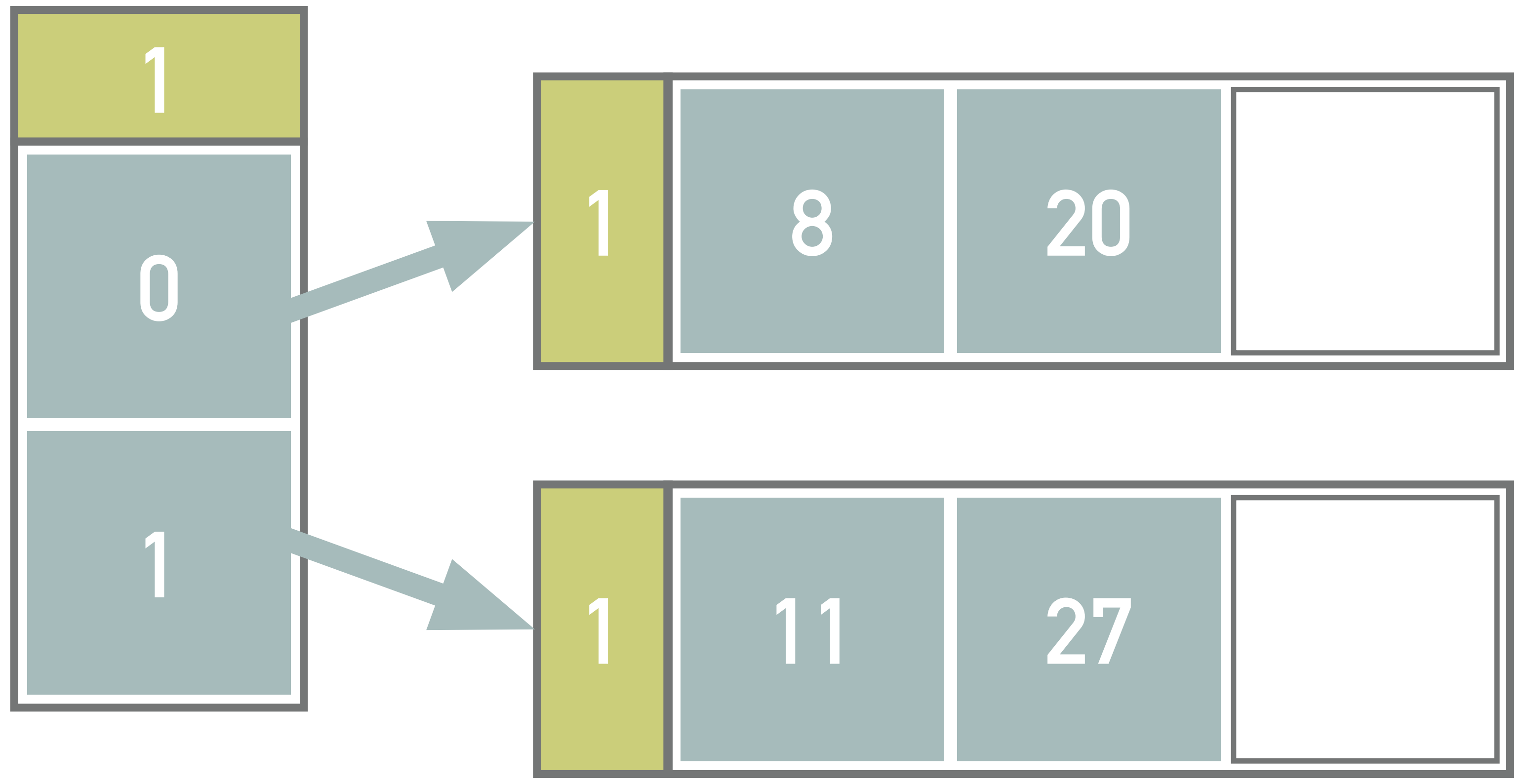
- $h(11_{10}) = 1011_2$
  - Record with key 11 is stored to the same page using entry 1

- $h(8_{10}) = 1000_2$
  - Record with key 8 is inserted into the same page using entry 0

EXAMPLE 2: FAGIN – SPLITTING A PAGE

- Insert record with key 27 into the structure from example 1
- $h(27_{10}) = 11011_2$
- Page is overflown
  - The local value d_L of the page is less than the global value d_G of the directory
  - Therefore we can split the page into two new pages and increment d_L values of both the pages
- Finally, we reinsert the records previously allocated into the page being split
  - After the reinsert, the even keys are stored in the page referenced from the zero-th directory entry while the odd records are referenced from the first entry





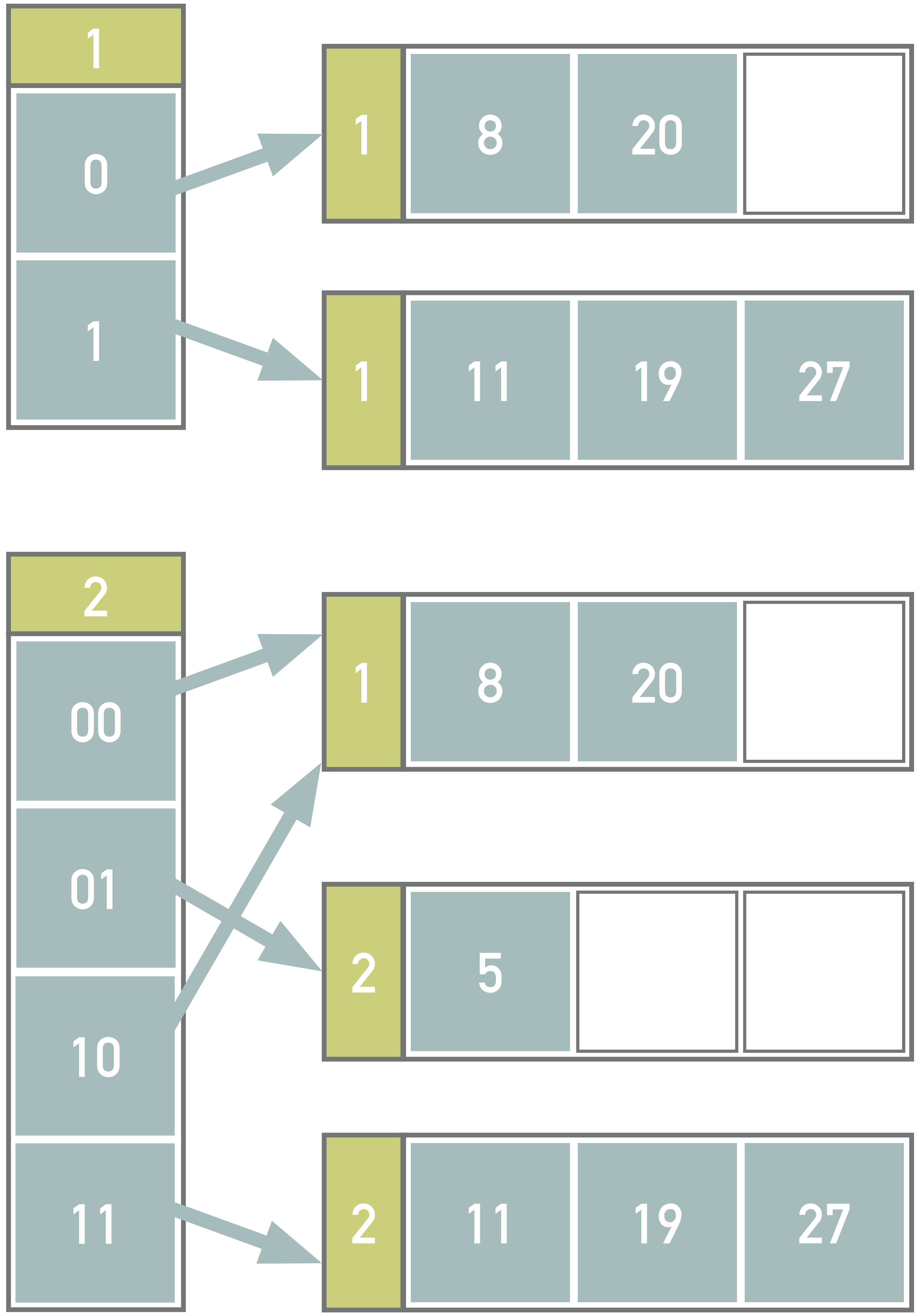
EXAMPLE 3: FAGIN – EXPANDING THE DIRECTORY

- Insert records with keys 19 and 5 into the structure from example 2

- $h(19_{10}) = 10011_2$
  - After inserting record with key 19, a page is filled

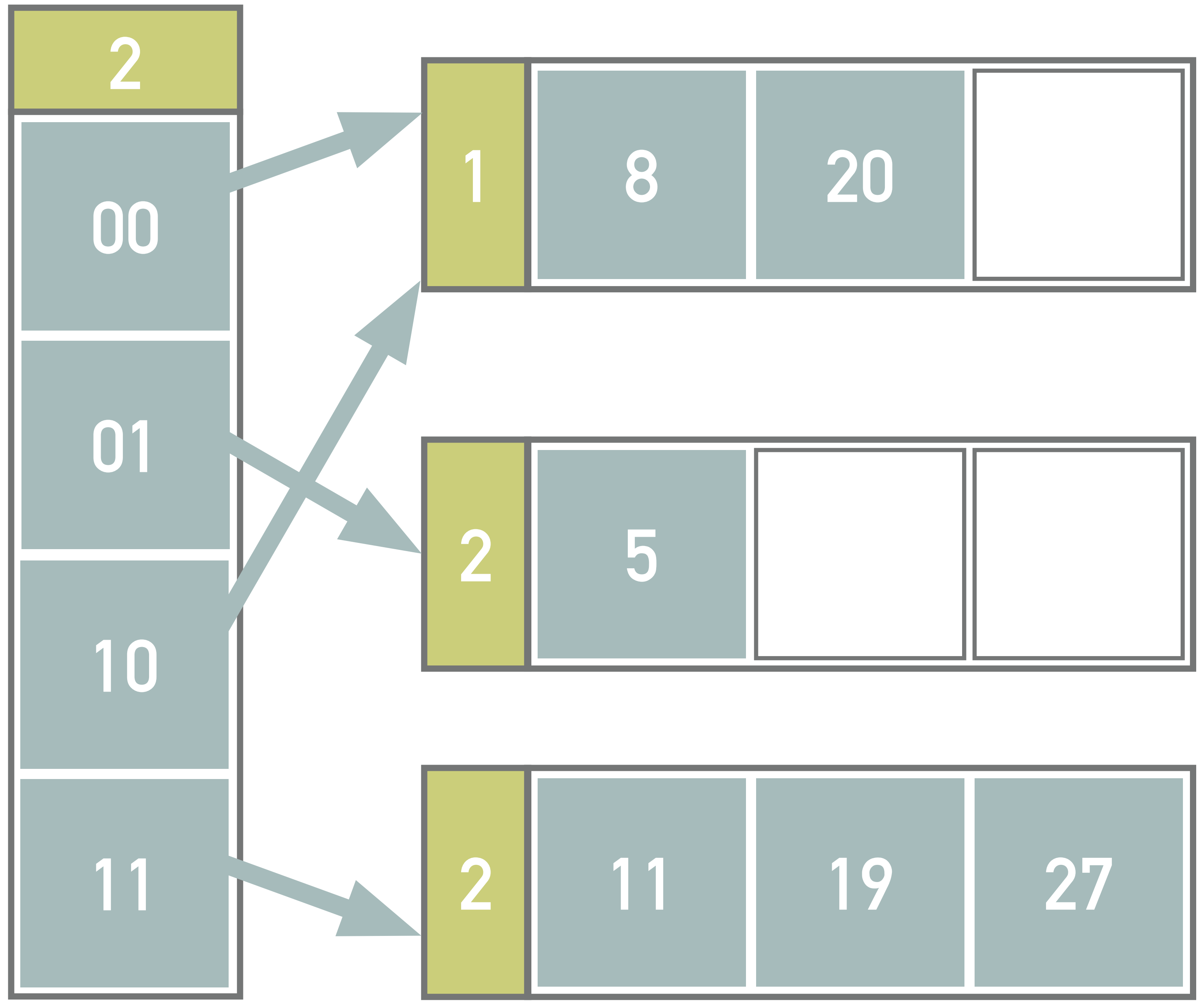
- $h(5_{10}) = 101_2$
  - The insert of the record having key 5 causes:
    - Expanding the directory, i.e., $d_L = d_G$
    - Splitting of the second page, i.e., $d_L = 2$
    - Reinserting of records with keys 5, 11, 19 and 27

- $h(11_{10}) = 1011_2$
- $h(19_{10}) = 10011_2$
- $h(27_{10}) = 11011_2$



EXERCISE 1

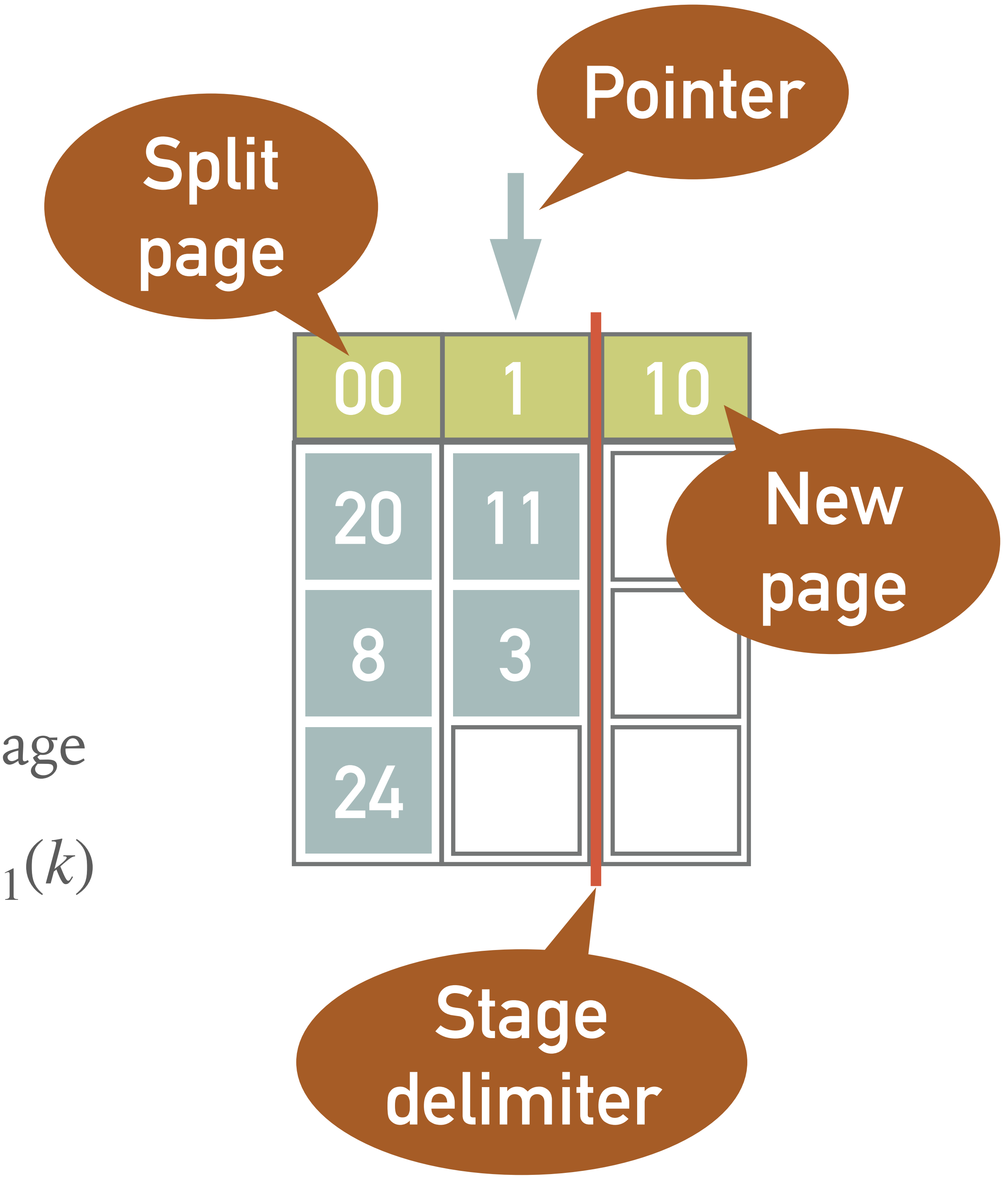
- Insert records with keys 24 and 32
- Note all the computations and illustrate the solution



LITWIN

- Directory-less scheme that avoids *exponentially increasing size* of the directory, but we need a continuous space in the secondary memory
    - Addition of a single page after *pre-defined condition*

- The primary file is *linearly expanded* with time (stages), i.e., adding one page after another
    - *Stage* d starts with $s = 2^d$ pages and ends when the number reaches $s = 2^{d+1}$ (i.e., stage $d + 1$ begins)
- During the stage, a *split pointer* $p \in \{0,\ldots,2^d - 1\}$ identifies the next page to be split
    - At the beginning of stage d, the pointer points to page 0
    - After every split operation, the *pointer is incremented* by 1, or moved to the start when we enter a new stage
    - Records from page p (and overflow pages) will be distributed between *split pages* p and $p + s$ using $h_{d+1}(k)$
    - If a page overflows before its time to split, *overflow page* will be utilised



- At each stage, we have two types of hash functions
    - $h_d(k)$ for pages not yet split, i.e., the least significant d bits of the hashed value $h(k)$ are used
    - $h_{d+1}(k)$ for the already split pages

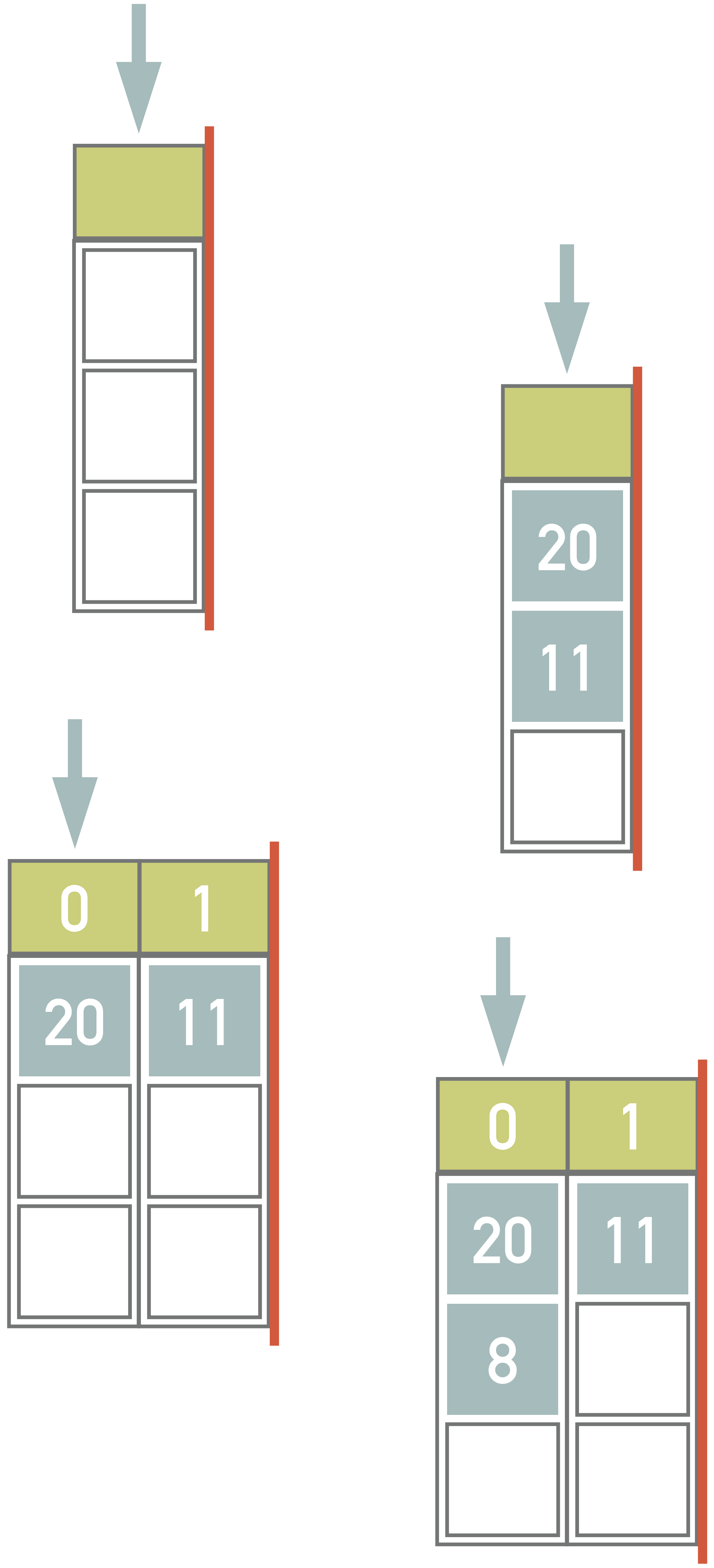
EXAMPLE 4: LITWIN

- Insert records with keys 20, 11 and 8 into an empty primary file
    - I.e., start the stage $d = 0$ with one page (capacity 3 records), $h(k) = k$, $p = 0$
    - Pre-defined condition: Splitting occurs after 2 inserts

- The records with key 20 and 11 are inserted into the 0-th page disregarding the value of the key
    - $d = 0$ bits of the keys are used at this point

- We have inserted 2 keys, therefore splitting occurs (a new page is created)
    - The records from 0-th page are redistributed using the least significant bit of the hashed key
        - $h(20_{10}) = 10100_2$
        - $h(11_{10}) = 1011_2$
    - Because $p = 2^1$ is reached, the stage changes to $d = 1$, $p = 0$

- Now, we use $d = 1$ bit for not yet split pages and $d + 1$ bits for split pages
    - The record with key 8 is inserted into the page 0 using the least significant bit
        - $h(8_{10}) = 1000_2$



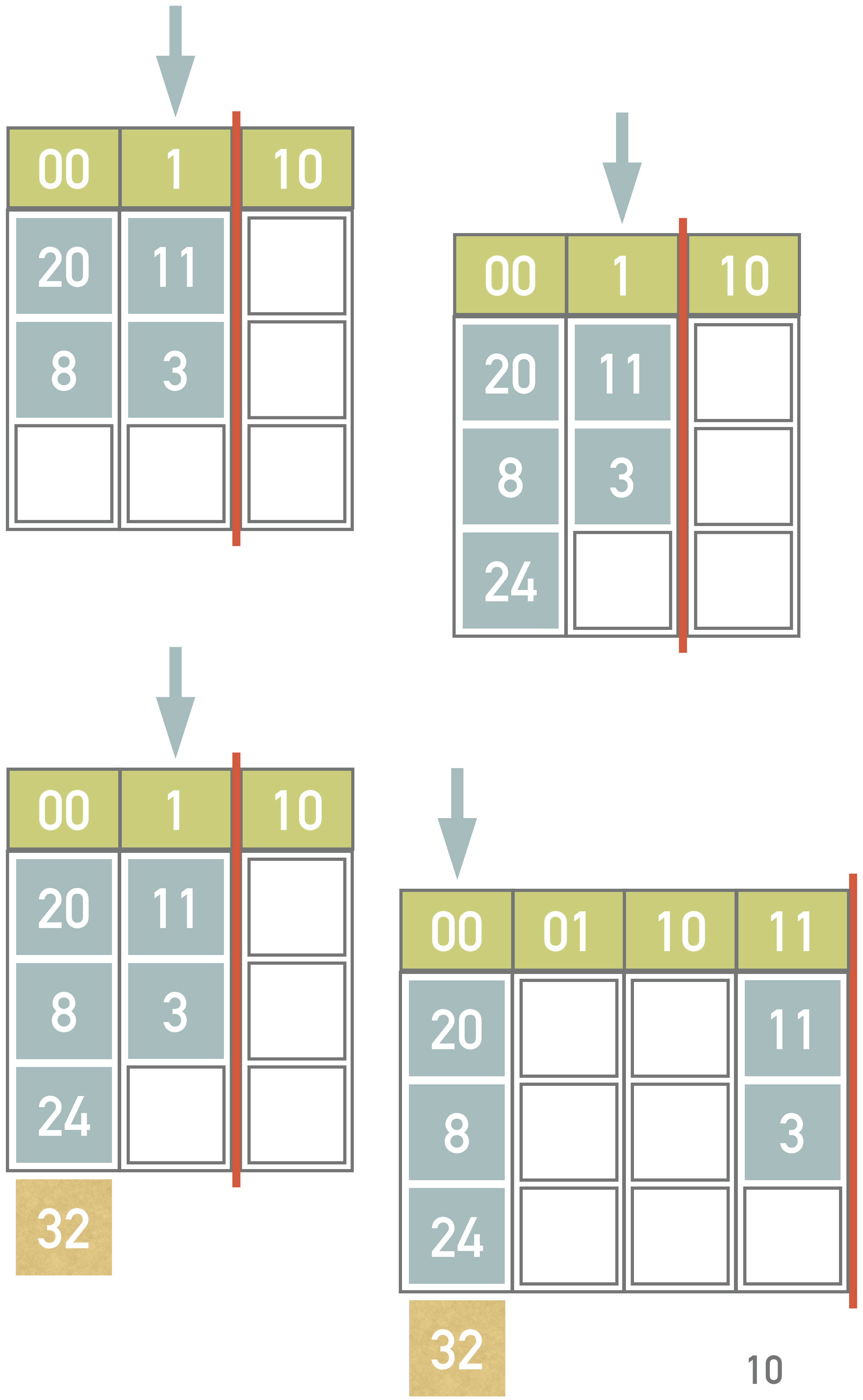
EXAMPLE 5: LITWIN

- Insert records with keys 3, 24 and 32 into the structure from example 4

- A record with key 3 will now be inserted into page 1
  - $h(3_{10}) = 11_2$

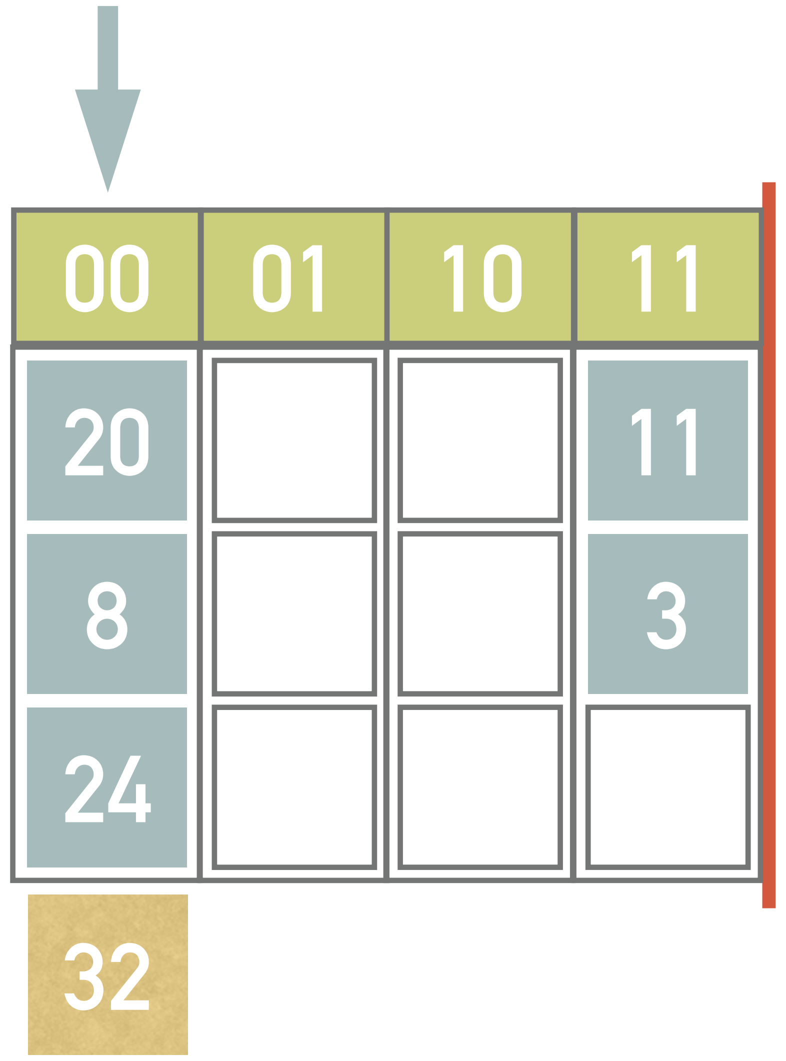
- We have already inserted 2 records in the stage $d = 1$, therefore page $p = 0$ is split into pages $p_0 = 00$, $p_1 = 10$

- Next, we will insert a pair of records with keys 24 and 32
  - $h(24_{10}) = 11000_2$
    - Because $h_1(11000_2) = 0$ and $0 < p$, it is necessary to address the keys using 2 least significant bits, i.e., $h_1(100000_2) = 0$, and the key belongs in the page 00
  - $h(32_{10}) = 100000_2$
    - The key 32 belongs to the same page, but that is already filled and thus overflows
  - Finally, the page 1 is split
    - Since the number of pages reaches $s = 2^{1+1} = 4$, the second stage is initiated, i.e., $d = 2$, $p = 0$



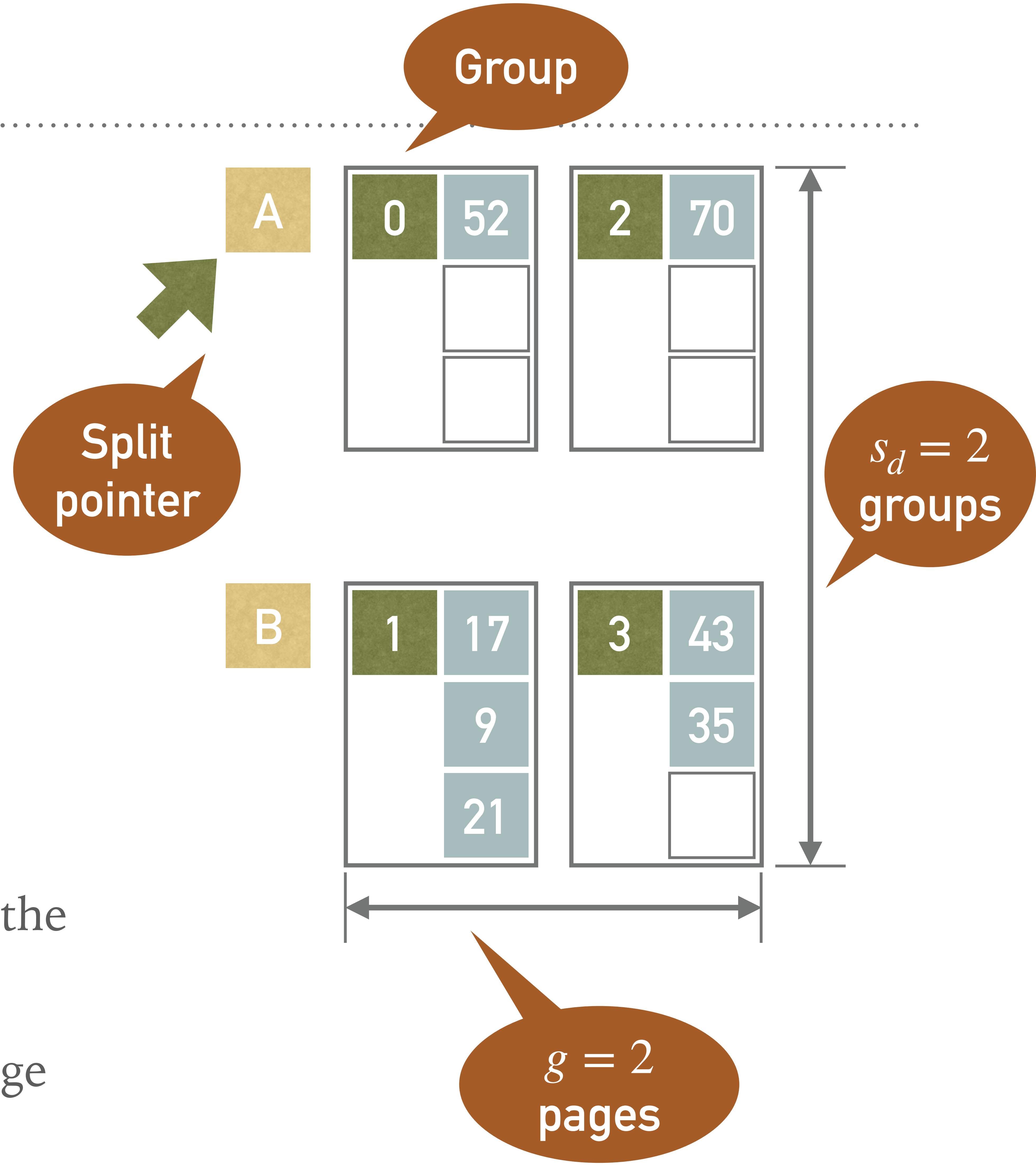
EXERCISE 2

- Insert records with keys 27, 19, 10, and 5 into the following structure
  - I.e., start the stage $d = 2$ with $s = 4$ pages (capacity 3 records), $h(k) = k$, $p = 0$
  - Pre-defined condition: Splitting occurs after 2 inserts
- Note all the computations and illustrate the solution



LHPE-RL

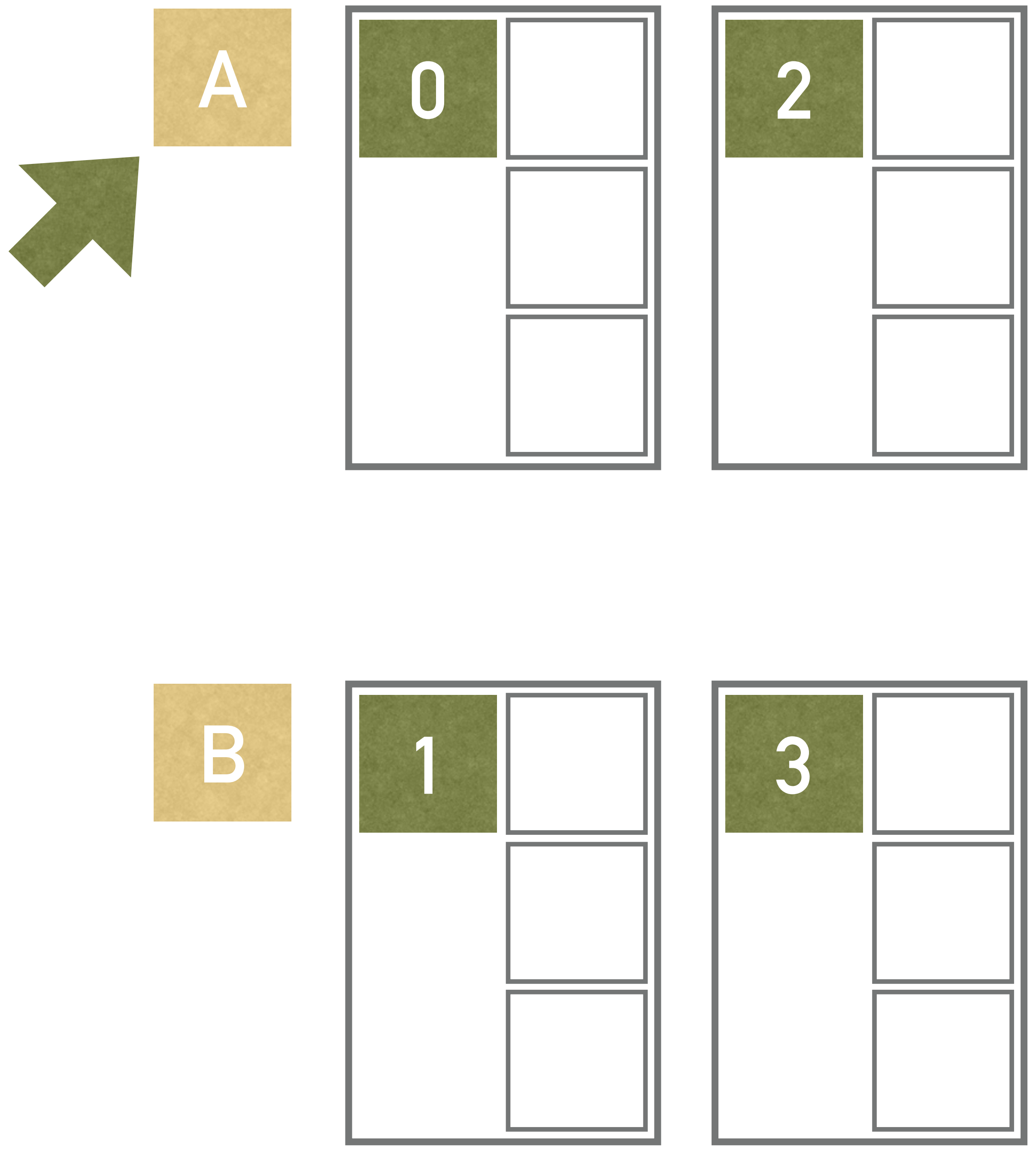
- Simplified version of LHPE
- At the *stage* d, the primary file consists of p_d *pages*
  - Each page has *capacity* b
  - Pages are grouped into $s_d = p_d \div g$ groups
    - Each *group* has g pages
- When a *predefined condition* is met (e.g., after L insertions), a *new page* is inserted at the end of the primary file and records in pages in the group pointed to by the *split pointer* are redistributed between these pages and the new page (being the new member of the group)
- When the last page is redistributed, the file is (virtually) *reorganized* (stage $d+1$) so that all the pages are again sorted into $s_{d+1} = p_{d+1} \div g$ pages
  - $p_{d+1} = \lceil s_d * (g+1) \div g \rceil * g$



EXAMPLE 6: LHPE-RL

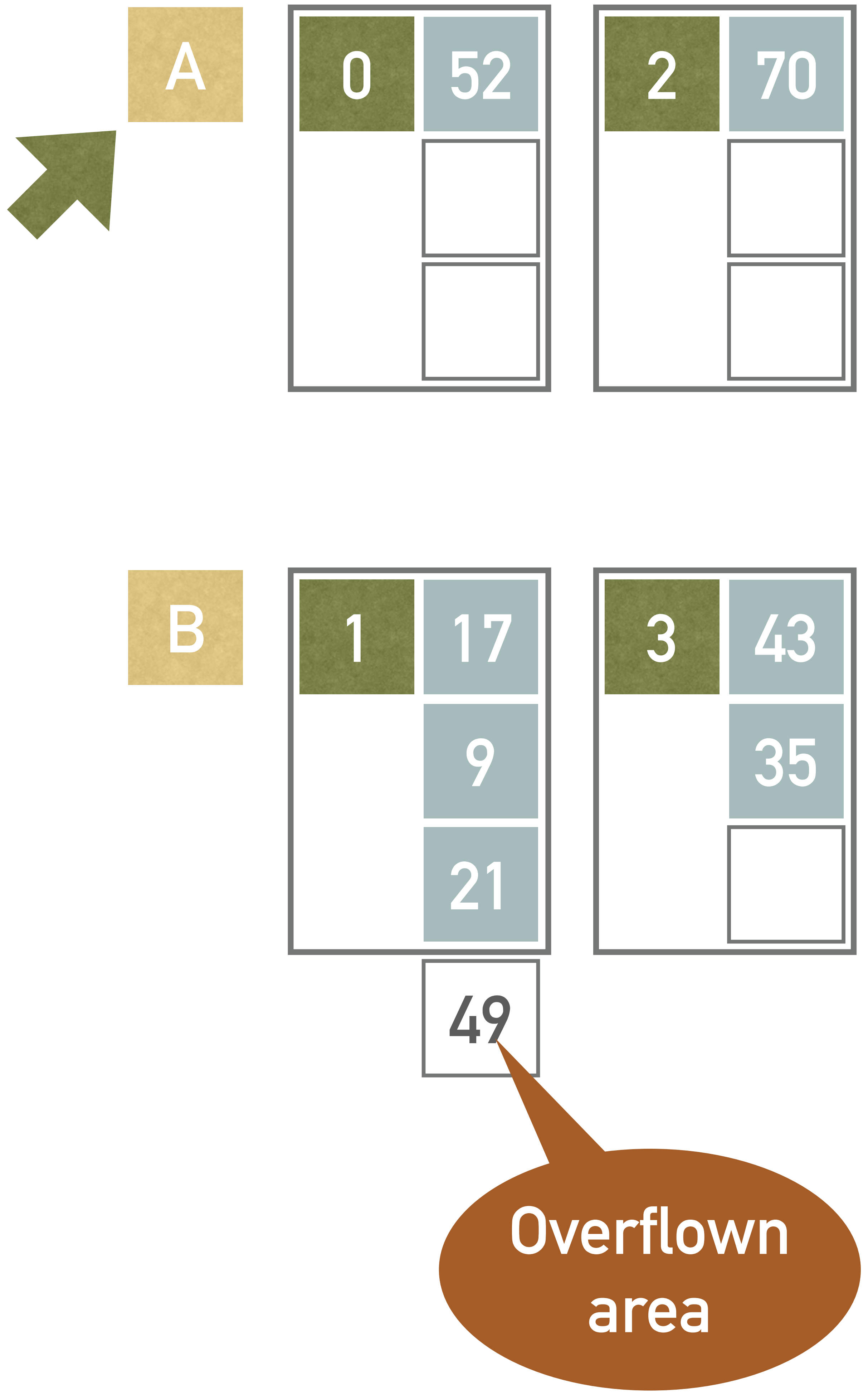
- Insert records with keys 17, 9, 43, 21, 49, 35, 70, 52, 40, 13, 5, 80 into the following empty structure
  - Stage $d = 0$
  - The initial number of groups $s_0 = 2$
  - Page capacity $b = 3$

- Hash functions
  - $h_0(k) = k \mod 4$
    - Determines into which of 4 initial pages a record is inserted at the beginning
  - $h_1(k) = k \mod 3$
    - Determines where the records are inserted when a group splits for the first time
  - $h_2(k) = (k \div 3) \mod 3$
    - Determines where the records are inserted when a group splits for the second time
- We are going to split regularly after two inserts, i.e., $L = 2$
  - We have $n = s \times g = 4$ pages, thus the first split happens after insertion of $n \times L = 8$ records



EXAMPLE 6: LHPE-RL

- Inserts of the first 8 keys, i.e., 17, 9, 43, 21, 49, 35, 70, 52, are not interesting since they are inserted where the h_0 function says
  - $h_0(17) = 17 \mod 4 = 1$
  - $h_0(9) = 9 \mod 4 = 1$
  - $h_0(43) = 43 \mod 4 = 3$
  - $h_0(21) = 21 \mod 4 = 1$
  - $h_0(49) = 49 \mod 4 = 1$
  - $h_0(35) = 35 \mod 4 = 3$
  - $h_0(70) = 70 \mod 4 = 2$
  - $h_0(52) = 52 \mod 4 = 0$

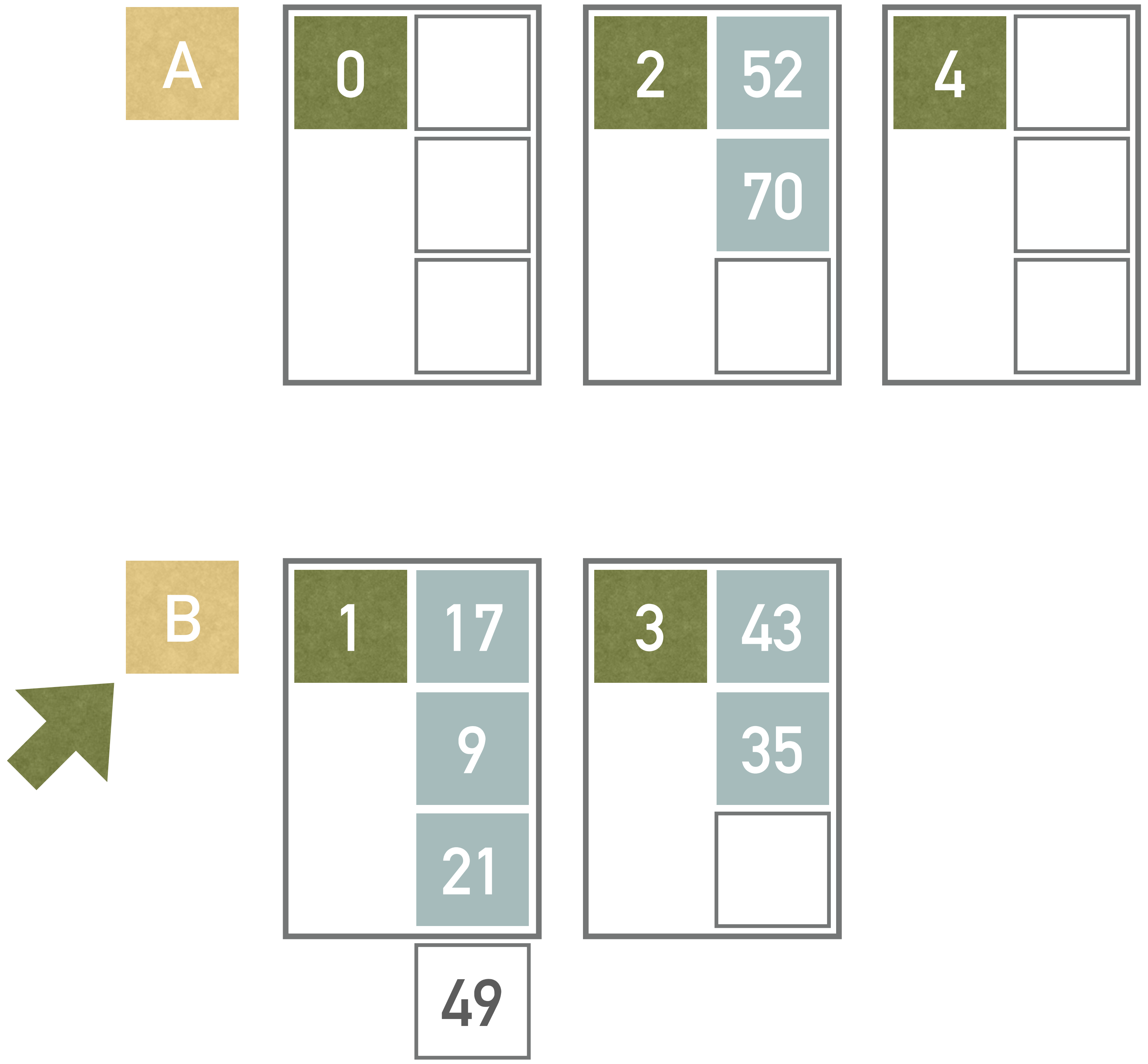


- The only problem is with key 49 which is assigned to a (already full) page 1

EXAMPLE 6: LHPE-RL

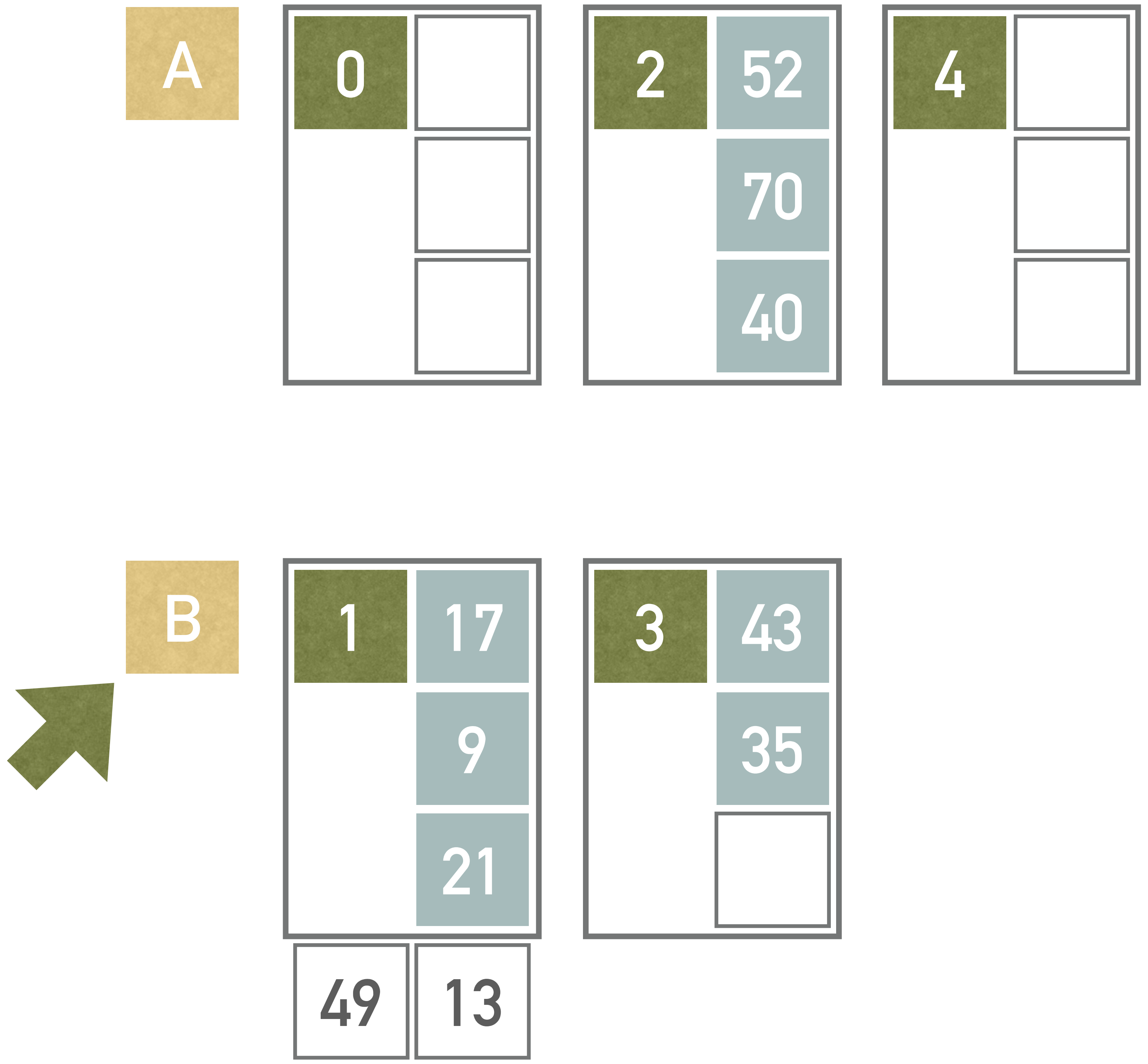
- We have inserted 8 keys so we have to split the group pointed by the split pointer, i.e., the group A having pages 0, and 2
  - Page 4 is added into the group A
  - Function $h_1(k)$ is applied in order to redistribute keys in the group A
    - $h_1(k)$ returns the index of a page in a group A, i.e., $h_1(k) = 0$ for the page 0, $h_1(k) = 1$ for the page 2, $h_1(k) = 2$ for the page 4
    - $h_1(52) = 52 \mod 3 = 1$, therefore key 52 goes into the page 2
    - $h_1(70) = 70 \mod 3 = 1$, therefore the key 70 goes into the page 2
  - Split pointer is incremented

- The key in the overflow area, i.e., 49, does not belong neither to page 0 nor to page 2, and thus stays where it is



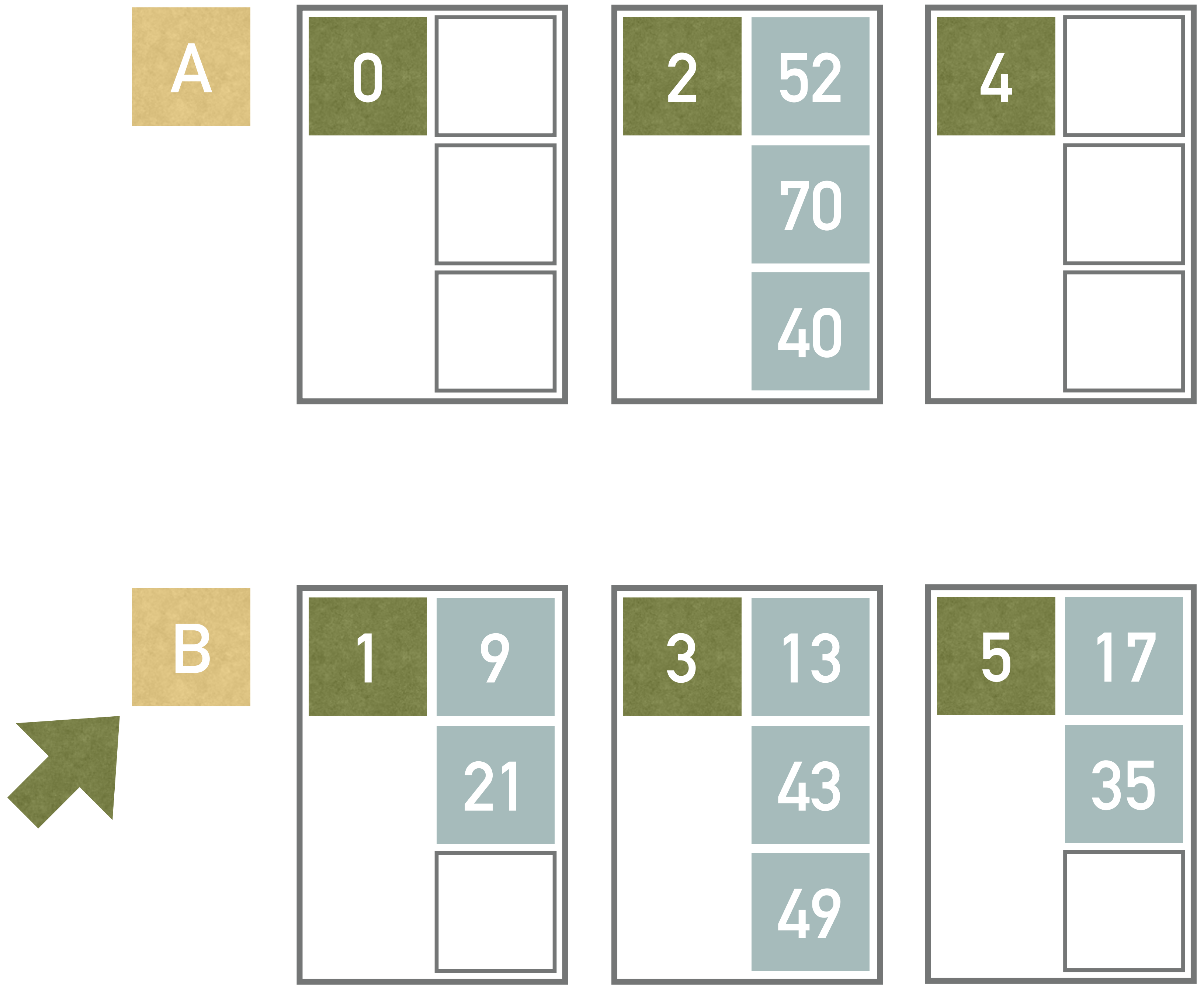
EXAMPLE 6: LHPE-RL

- Next, we insert record with key 40
  - $h_0(40) = 40 \mod 4 = 0$
  - Based on the function h_0, the record with key 40 should be assigned to the page 0 but this page has already been split
  - Therefore we need to use h_1 which sends it into the second page in the group A (page 2)
    - $h_1(40) = 40 \mod 3 = 1$
- Next, we insert record with key 13
  - $h_0(13) = 13 \mod 4 = 1$
  - Based on the function h_0, the record with key 13 belongs to the page 1, which has not been split yet
    - No need to use h_1
  - The page 1 is full, therefore the overflow area is used



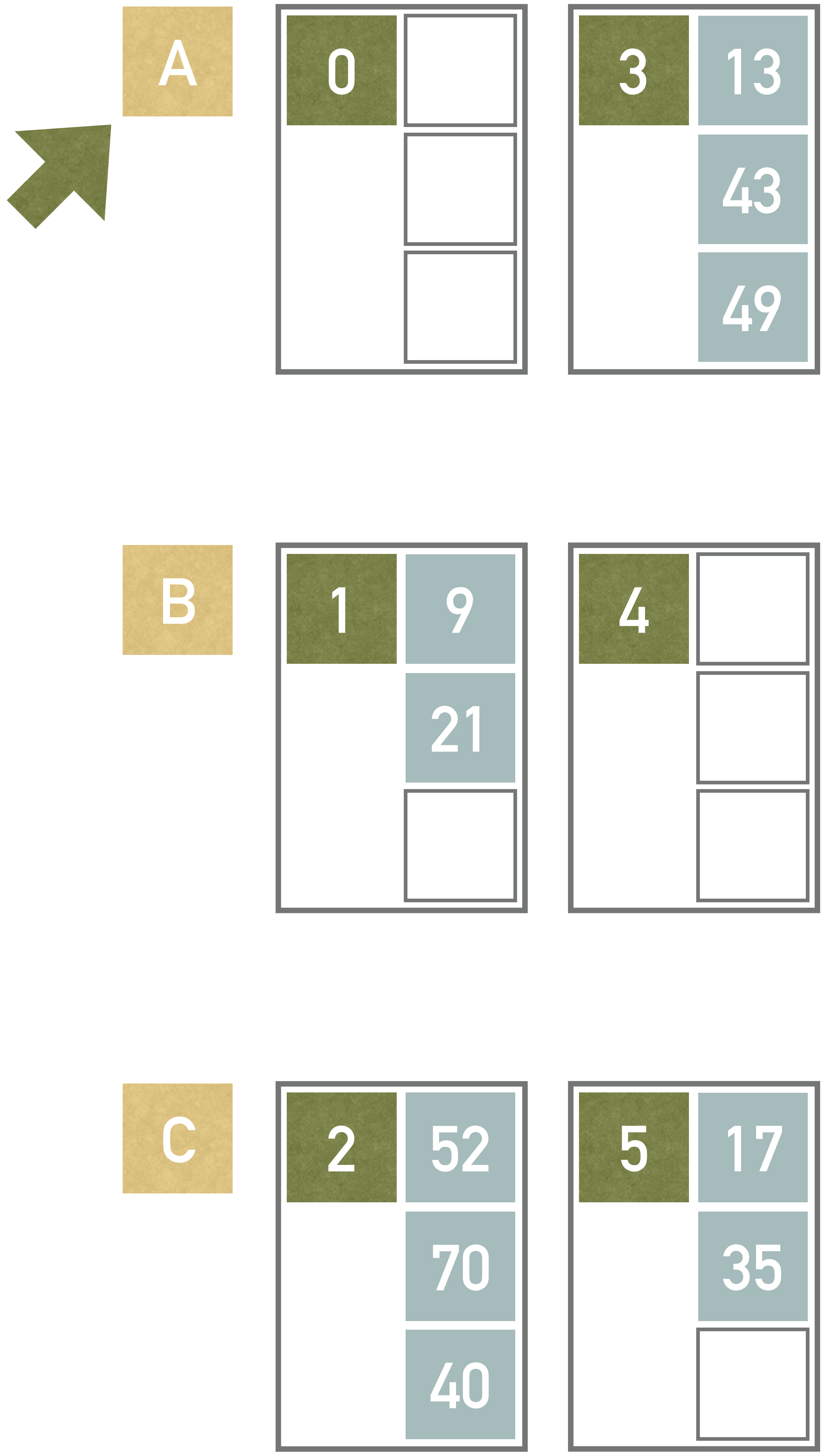
EXAMPLE 6: LHPE-RL

- Once again, we have to split the group (we have already inserted $L = 2$ records)
  - Split pointer points to the group B, i.e., pages 1 and 3 will be split
  - Page 5 is added
  - Function $h_1(k)$ will be applied in order to redistribute keys in the group B
    - $h_1(17) = 17 \mod 3 = 2$, therefore goes to the page 5
    - $h_1(9) = 9 \mod 3 = 0$, therefore goes to the page 1
    - $h_1(21) = 21 \mod 3 = 0$, therefore goes to the page 1
    - $h_1(43) = 43 \mod 3 = 1$, therefore goes to the page 3
    - $h_1(35) = 35 \mod 3 = 2$, therefore goes to the page 5
    - $h_1(49) = 49 \mod 3 = 1$, therefore goes to the page 3
    - $h_1(13) = 13 \mod 3 = 1$, therefore goes to the page 3



EXAMPLE 6: LHPE-RL

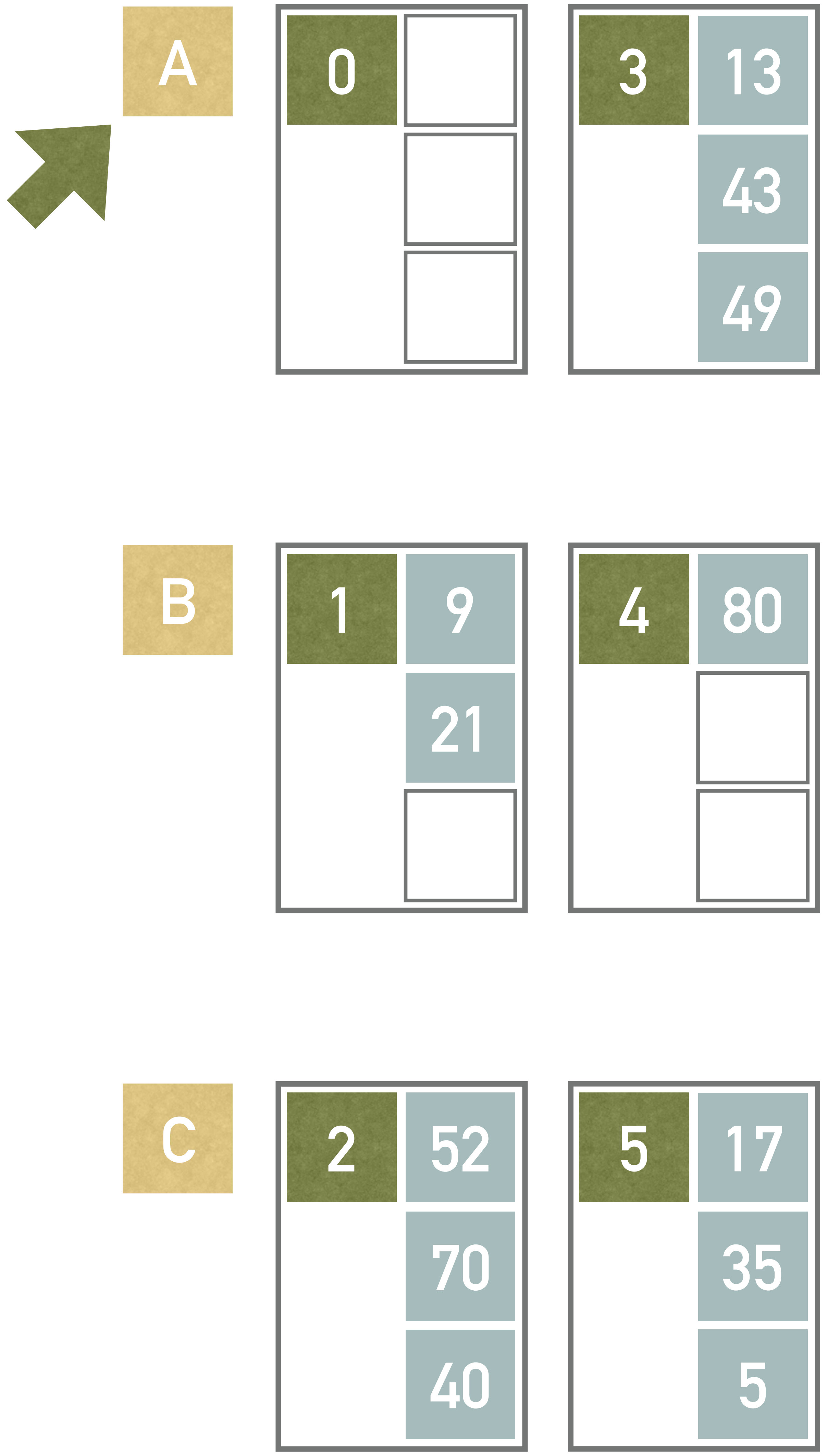
- Having all the groups processed (by split operation), the end of the stage $d = 0$ occurs
  - We will reorganize the file into 3 groups, each having two pages
- The reorganization is only virtual
  - The page numbers are kept, we just think of the pages differently



EXAMPLE 6: LHPE-RL

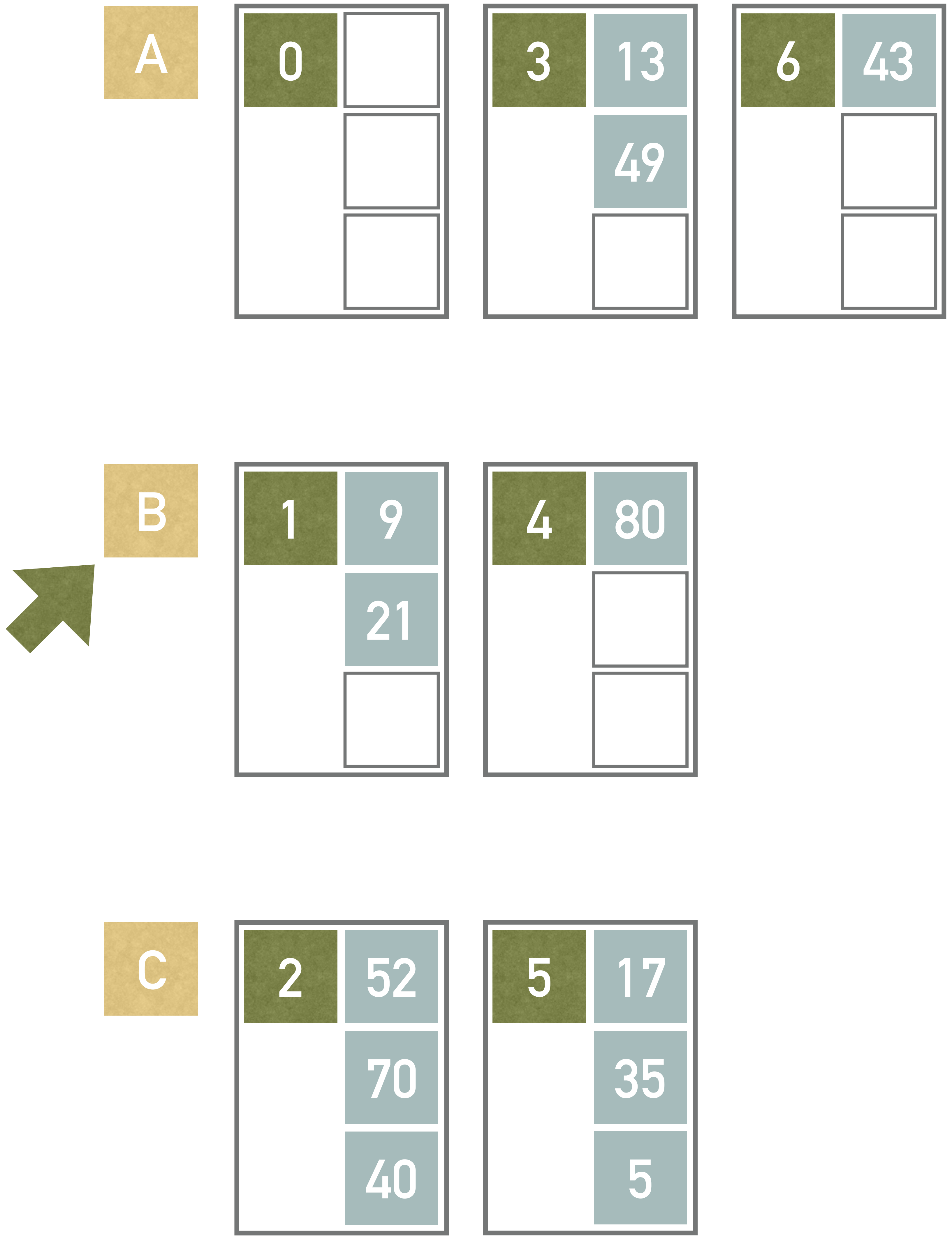
- Now, we insert record with key 5
  - $h_0(5) = 5 \mod 4 = 1$
  - Based on the function h_0, this record belongs to the page 1, but this has been split once
    - Therefore we have to use h_1
    - $h_1(5) = 5 \mod 3 = 2$ (note that redistribution is only virtual)
  - The record comes into page 5

- Next, we insert record with key 80
  - $h_0(80) = 80 \mod 4 = 0$
  - Based on the function h_0, this record belongs to the page 0, but this has been split once
    - Therefore we have to use h_1
    - $h_1(80) = 80 \mod 3 = 2$ (note that redistribution is only virtual)



EXAMPLE 6: LHPE-RL

- Having inserted additional $L = 2$ records, we must split once again
  - The split pointer points to the group A, i.e., pages 0 and 3
  - Page 6 is added into the group A
  - Function $h_2(k)$ is applied in order to redistribute keys in the group A
    - $h_2(k)$ returns the index of a page in a group A, i.e., $h_2(k) = 0$ for the page 0, $h_2(k) = 1$ for the page 3, $h_2(k) = 2$ for the page 6
    - $h_2(43) = (43 \div 3) \mod 3 = 2$ -> page 6
    - $h_2(49) = (49 \div 3) \mod 3 = 1$ -> page 3
    - $h_2(13) = (13 \div 3) \mod 3 = 1$ -> page 3
- Finally, split pointer is incremented



EXERCISE 3

- Insert records with keys 37 into the structure from example 6 (see the picture)
  - Stage $d = 1$
  - Page capacity $b = 3$
  - Predefined condition $L = 2$
  - Hash functions:
    - $h_0(k) = k \mod 4$
    - $h_1(k) = k \mod 3$
    - $h_2(k) = (k \div 3) \mod 3$
- Note all the computations and illustrate the solution