

# PRINCIPLES OF DATA ORGANISATION

File Organization



# MOTIVATION

- ⌘ How to organize data?
- ⌘ Implementation



# HEAP FILE

- ⌘ Note: not heap!!
- ⌘ Variable-length records ~ file is not homogeneous
  - ⌘ E.g., a log file
- ⌘ Data not sorted in any way, a record placed always at the end of the file
- ⌘ Usually used along with another supporting structure
- ⌘ Insert :  $O(1)$ 
  - ⌘ Fetch the last block (keep in memory) in the file and append the new record
- ⌘ Find :  $O(N / b)$ 
  - ⌘  $N$  – number of records
  - ⌘  $b$  – average blocking factor
    - ⌘  $b = \lfloor B/R \rfloor$ 
      - ⌘  $B$  = block size
      - ⌘  $R$  = record size
- ⌘ Whole file needs to be scanned



# UNSORTED SEQUENTIAL FILE

- Fixed-size record
- Data not sorted in any way
  - Heap file
- Suitable when data are collected without any relationship to other data
  - We can query for a record using its index

- Insert :  $O(1)$
- Find :  $O(N / b)$

Block	Name	Department	...
0	Galvin Janice	Purchasing	
	Walters Rob	Marketing	
	Brown Kevin	Marketing	
1	Walters Rob	Development	
	Duffy Terri	Research	
	Brown Kevin	PR	
2	Duffy Terri	Development	
	Walters David	Production	
	Brown Kevin	Purchasing	
3	Matthew Gigi	Purchasing	
	Walters Rob	PR	
	...	...	



# SORTED SEQUENTIAL FILE

- Fixed size record
- Records sorted in the file according to the primary search key
  - According to only one – the most often searched

## Fetch

- Sequential scan
  - $O(N/b)$
- Binary search
  - Direct-access medium
  - $O(\log(N/b))$
- Range query
  - Find start and read  $k$  records

Block	Name	Department	...
0	Brown Kevin	PR	
	Brown Kevin	Purchasing	
1	Brown Kevin	Marketing	
	Duffy Terri	Development	
2	Duffy Terri	Research	
	Galvin Janice	Purchasing	
3	Matthew Gigi	Purchasing	
	Walters David	Production	
4	Walters Rob	Marketing	
	Walters Rob	Development	
5	Walters Rob	PR	
	...	...	



# SORTED SEQUENTIAL FILE — MODIFICATION

## Insert

- ↳ Inserting a new record is costly
  - ↳ All the following records would have to be shifted
- ↳ Auxiliary file/blocks called **overflow file/bucket** need to be established where the new records are inserted
  - ↳ Outside the primary file
- ↳ The (main) file is **periodically** reorganized

## Update

- ↳ Simple if the update does not include the primary search key
  - ↳ If so, it is delete and insert

## Delete

- ↳ Deleted records are not directly removed
  - ↳ Reorganization would have to take place
- ↳ A bit designating deleted records is set
- ↳ Deleted records are removed during **periodical reorganization**

What if we want to access the data using various attributes?



# INDEX

## ↳ Motivation:

- ↳ Can we do it better than sorting the data?

- ↳ Yes

- ↳ Recall binary tree, a-b tree, ...

↳ An **index** is an auxiliary structure for a data file that consists of a specifically arranged structure containing key-pointer pairs

- E.g., name-pointer to the block with the record

## ↳ Storage of the index

- ✦ Main memory

- ✦ Cashed

- ✦ Secondary memory

- ✦ Accessing index must also be taken into account when computing the find/fetch time

- ✦ In real use: blocking factor of the index >> blocking factor of the primary file



# INDEXED SORTED SEQUENTIAL FILE

Fixed size records

Structure

Primary file/area

Data sorted according to the primary search key

Index/secondary file/area

Typically hierarchical

Overflow file/area

Data can be accessed either using the index or sequentially

$b = 2$

Primary file

$b = 5$

Index file  
1. (base) level

Index file  
2. (top) level

Block	Name	
8	Brown Kevin	6
	Walters Rob	7

Block	Name	
6	Brown Kevin	0
	Clinard Stephnie	1
	Duffy Terri	2
	Leavy Shirleen	3
	Peagler David	4
	Walters Rob	5
	...	

Block	Name	...
0	Brown Kevin	
	Berkman Doloris	
1	Clinard Stephnie	
	Coolidge Emily	
2	Duffy Terri	
	Galvin Janice	
3	Leavy Shirleen	
	Matthew Gigi	
4	Peagler David	
	Shackelford Elsie	
5	Walters Rob	

# INDEXED SEQUENTIAL FILE — FETCH

Searching for a specific value (query key)

- ④ Check the top level of the index and identify a key-value pair with the highest value lower than the query key
- ④ Fetch the block referenced by the value
- ④ Repeat the previous steps with lower index levels until a primary file block is reached
  - ④ Fetch time depends on the **height of the tree**
  - ④ Each level = disc access
- ④ Search the primary file block for the specified key

Searching for a range of values

- ④ Search for the lower bound key of the interval
- ④ Sequentially scan the blocks of the primary file until the record corresponding to the upper bound key is found



# INDEXED SEQUENTIAL FILE – FETCH

Search for Galvin Janice (G)

Search for Galvin Janice – Walters Rob (G – W)

$b = 2$

Primary file

Block	Name	...
0	Brown Kevin	
	Berkman Doloris	
1	Clinard Stephnie	
	Coolidge Emily	
2	Duffy Terri	
	Galvin Janice	
3	Leavy Shirleen	
	Matthew Gigi	
4	Peagler David	
	Shackelford Elsie	
5	Walters Rob	

$b = 5$

Index file  
1. (base) level

Block	Name	
6	Brown Kevin	0
	Clinard Stephnie	1
	Duffy Terri	2
	Leavy Shirleen	3
	Peagler David	4
7	Walters Rob	5
	...	

Index file  
2. (top) level

Block	Name	
8	Brown Kevin	6
	Walters Rob	7

# INDEXED SEQUENTIAL FILE — INSERT

- ↳ When an index is created, index nodes are fixed and do not change during modifications of the primary file
  - ✦ Index structure is **static**
    - ↳ Later we will see that it does not have to be
- ↳ New records need to be stored in reserved areas (**pockets**) within the primary file
  - Long pockets decrease efficiency
- ↳ Overflown data are inserted into a new block (created dynamically) – **overflown block**
  - Outside the primary file
  - Buckets can be chained and therefore theoretically the ISF does not need to be rebuilt
    - ↳ But decrease performance
- ↳ Pointer to the overflow area
  - ✦ for each record in a block
    - More space
    - Shorter sequences in the overflow area
  - ✦ for each block



# INDEXED SEQUENTIAL FILE

## Pros

- ⌘ Fast access using primary search key
- ⌘ Shares pros of the sequential file

## Cons

- ⌘ Fast access only when using primary search key
- ⌘ Otherwise sequential scan
- ⌘ Problems with primary file when updating
- ⌘ Pockets slow down data access
- ⌘ Occasional reorganisation (also slow)



# INDEXED FILE ORGANIZATION

- ↳ Allows to search the file according to **different attributes** without the need to scan the whole file sequentially
- ↳ The primary file stays **unsorted** or is **sorted** according to one key only (primary index)
  - ↳ Sorted = we need to keep the ordering
    - ↳ If sorted using an artificial key, range queries are not common
  - ↳ Unsorted – e.g., heap with additional smart structure
- ↳ For each **query key** an index file can be built
  - one primary data file, multiple index files
- ↳ Basically corresponds to a standard database table
  - ↳ One table
  - ↳ Multiple indexes built over it (possibly of different types)



# INDEX

## Primary index

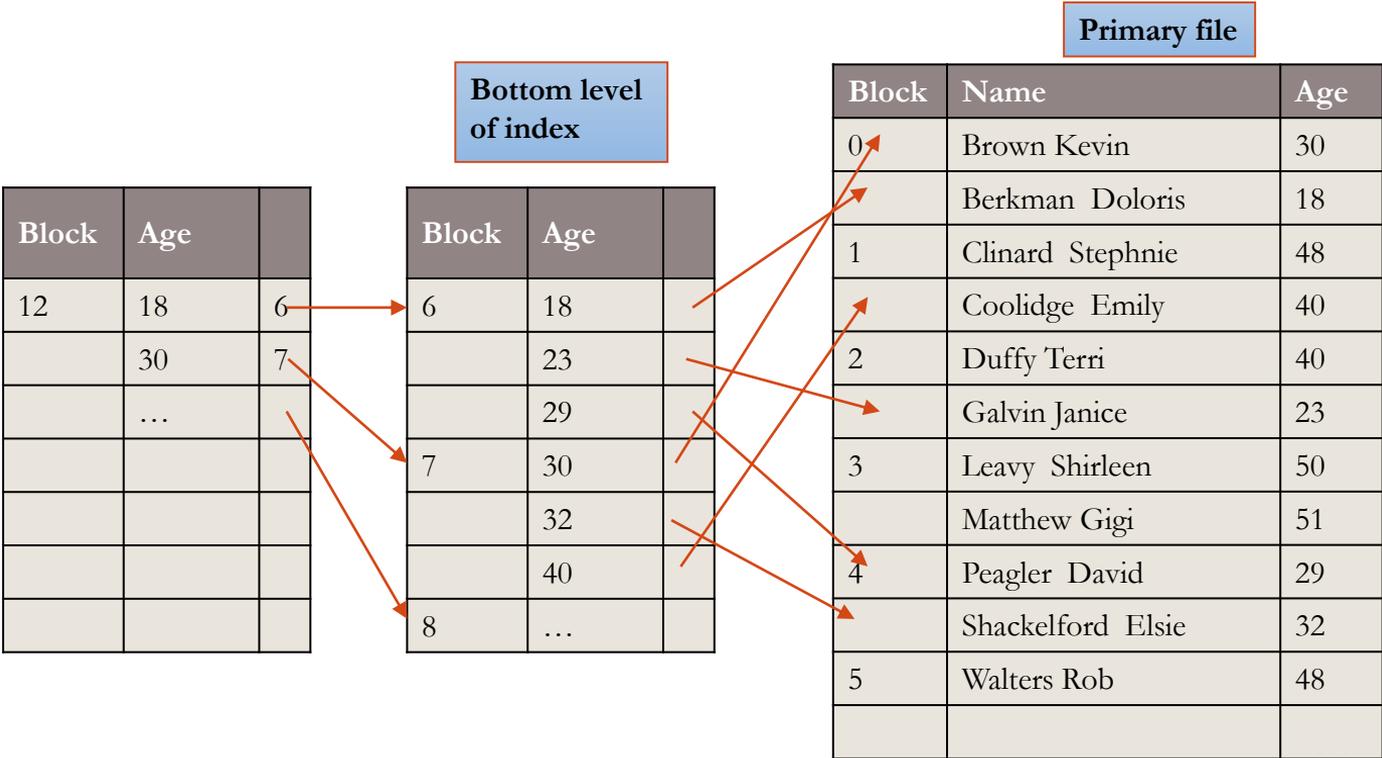
- ↳ Index over the attribute based on which the **records in the primary file are sorted**
  - ↳ Only one
- ↳ If the value of the primary attribute is modified, the file needs to be reorganised → should be relatively invariable
- ↳ Well-suited for range queries
- ↳ There does not have to be a primary index
- ↳ It is desirable to keep it in memory
  - ↳ Small keys (integer, not string)

## Secondary index

- ↳ There can be multiple secondary indexes
- ↳ We do not index blocks of the primary file, but a sorted list of indexed values (with pointers to the blocks with the data)
  - ↳ The bottom level of the index = we index records, not blocks
  - ↳ Next levels = we index blocks (with sorted records)
- ↳ Range queries for long ranges can be very expensive



# SECONDARY INDEX



# INDEX

## Direct index

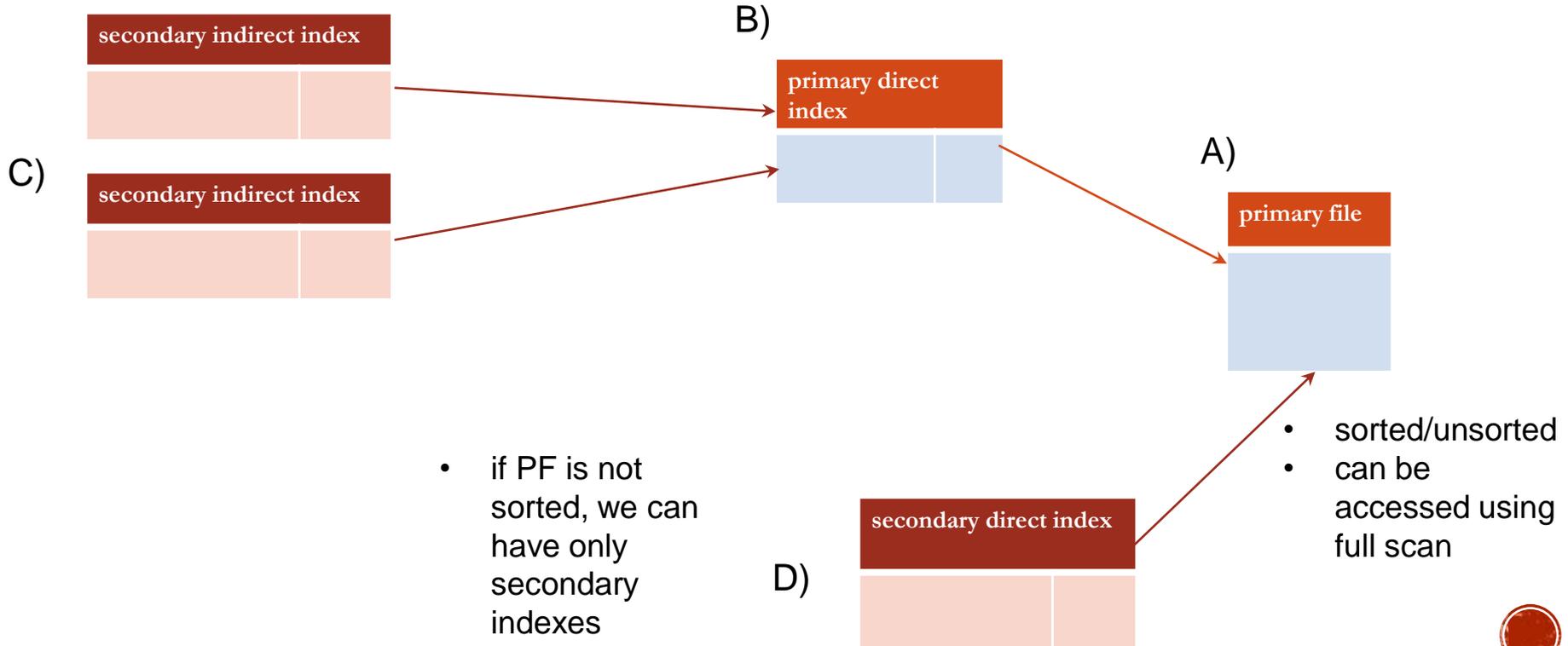
- ↳ Index is bound directly to records
  - ↳ Pointers to the primary data file
- ↳ Primary file reorganization → modification of indexing structures

## Indirect index

- ↳ Contains **keys** of the data (which are in the primary index)
  - ↳ Not pointers to the primary file
- ↳ Accessing a record needs one more access to the primary index
- ↳ If the primary file is reorganised, the secondary indexes stay intact



# INDEX



# HASHED FILE ORGANIZATION

- Direct access with **one unique key**
- Use hash function to map records to pages/blocks addresses
- If the data can not fit into a page/block when inserting, an overflow strategy is employed
- Placement within the page is not specified
- When file is being reorganized, the pages are filled only to, e.g., 80%
  - To avoid overflow with next insert
  - The value depends on expected insert count



