Czech Technical University in Prague, Faculty of Information Technology

MIE-PDB: Advanced Database Systems

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Lecture 12

MapReduce, Apache Hadoop

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Lecture Outline

MapReduce

- Programming model and implementation
- Motivation, principles, details, ...

Apache Hadoop

- HDFS Hadoop Distributed File System
- MapReduce

Programming Models

What is a programming model?

- Abstraction of an underlying computer system
 - Describes a logical view of the provided functionality
 - Offers a public interface, resources or other constructs
 - Allows for the expression of algorithms and data structures
 - Conceals physical reality of the internal implementation
 - Allows us to work at a (much) higher level of abstraction
- The point is how the intended user thinks in order to solve their tasks and not necessarily how the system actually works

Programming Models

Examples

- Traditional von Neumann model
 - Architecture of a physical computer with several components such as a central processing unit (CPU), arithmetic-logic unit (ALU), processor registers, program counter, memory unit, etc.
 - Execution of a stream of instructions
- Java Virtual Machine (JVM)
- ...

Do not confuse programming models with

- Programming paradigms (procedural, functional, logic, modular, object-oriented, recursive, generic, data-driven, parallel, ...)
- Programming languages (Java, C++, ...)

Programming Models

Parallel Programming Models

Process interaction

Mechanisms of mutual communication of parallel processes

- Shared memory shared global address space, asynchronous read and write access, synchronization primitives
- Message passing
- Implicit interaction

Problem decomposition

Ways of problem decomposition into tasks executed in parallel

- Task parallelism
- Data parallelism independent tasks on <u>disjoint partitions of data</u>
- Implicit parallelism

MapReduce

MapReduce Framework

What is MapReduce?

- Programming model + implementation
- Developed by Google in 2008

Google:

A simple and powerful interface that enables automatic parallelization and distribution of large-scale computations, combined with an implementation of this interface that achieves high performance on large clusters of commodity PCs.

MapReduce Framework

A bit of history and motivation

Google PageRank problem (2003)

- How to rank tens of billions of web pages by their importance
 - ... <u>efficiently</u> in a reasonable amount of time
 - ... when data is scattered across thousands of computers
 - ... data files can be enormous (terabytes or more)
 - ... data files are updated only occasionally (just appended)
 - ... sending the data between compute nodes is expensive
 - ... hardware failures are rule rather than exception
- Centralized index structure was no longer sufficient
- Solution
 - Google File System a distributed file system
 - MapReduce a programming model

MapReduce Framework

MapReduce programming model

- Cluster of commodity personal computers (nodes)
 - Each running a host operating system, mutually interconnected within a network, communication based on IP addresses, ...
- Data is distributed among the nodes
- Tasks executed in parallel across the nodes

Classification

- Process interaction: message passing
- Problem decomposition: data parallelism

MapReduce Model

Basic Idea

Divide-and-conquer paradigm

- Map function
 - Breaks down a problem into sub-problems
 - Processes input data in order to generate a set of intermediate key-value pairs
- Reduce function
 - Receives and combines sub-solutions to solve the problem
 - Processes and possibly reduces intermediate values associated with the same intermediate key

And that's all!

MapReduce Model

Basic Idea

And that's all!

- It means...
 - We only need to implement Map and Reduce functions
 - Everything else such as
 - input data distribution,
 - scheduling of execution tasks,
 - monitoring of computation progress,
 - inter-machine communication,
 - handling of machine failures,
 - ..

is managed automatically by the framework!

MapReduce Model

A bit more formally...

Map function

- Input: an input key-value pair (input record)
- Output: a set of intermediate key-value pairs
 - Usually from a different domain
 - Keys do not have to be unique
- $(key, value) \rightarrow list of (key, value)$

Reduce function

- Input: an intermediate key + a set of (all) values for this key
- Output: a possibly smaller set of values for this key
 - From the same domain
- $(key, list of values) \rightarrow (key, list of values)$

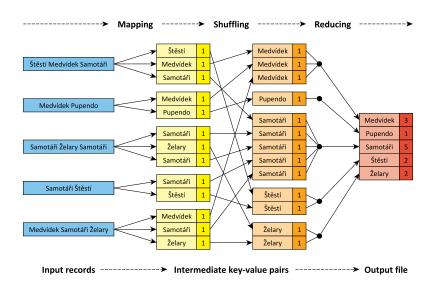
Example: Word Frequency

Implementation

```
/**
 * Map function
 * @param key Document identifier
 * @param value Document contents
 */
map(String key, String value) {
 foreach word w in value: emit(w, 1);
}
```

```
/**
 * Reduce function
 * @param key Particular word
 * @param values List of count values generated for this word
 */
reduce(String key, Iterator values) {
 int result = 0;
 foreach v in values: result += v;
 emit(key, result);
}
```

Logical Phases



Logical Phases

Mapping phase

- Map function is executed for each input record
- Intermediate key-value pairs are emitted

Shuffling phase

 Intermediate key-value pairs are grouped and sorted according to the keys

Reducing phase

- Reduce function is executed for each intermediate key
- Final output is generated

Framework Architecture

Master-slave architecture

- Master
 - Manages the entire execution of MapReduce jobs
 - Schedules individual Map / Reduce tasks to idle workers
- Slave (worker)
 - Accepts Map / Reduce tasks from the master

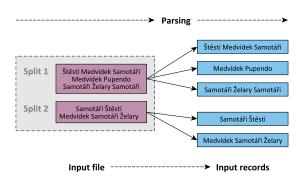
Input / output files

- Stored in the underlying distributed file system
- Actual contents of these files...
 - Divided into smaller splits
 - Physically stored by individual slaves

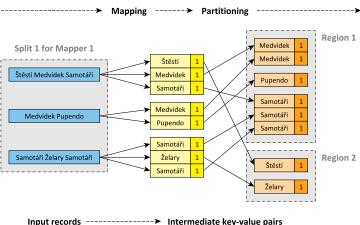
Input Parsing

Parsing phase

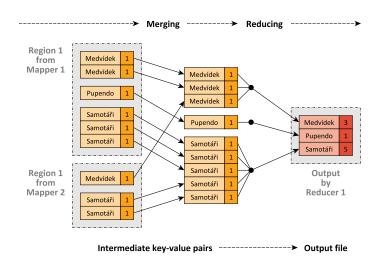
 Each split is parsed so that input records are retrieved (i.e. input key-value pairs are obtained)



Mapping Phase



Reducing Phase



Execution Functions

Input reader

Parses a given input split and prepares input records

Map function

Partition function

• Determines a particular Reducer for a given intermediate key

Compare function

Mutually compares two intermediate keys

Combine function

Reduce function

Output writer

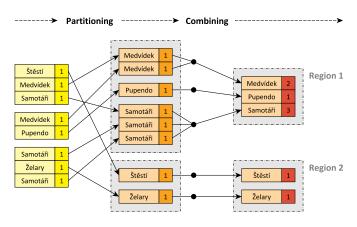
Writes the output of a given Reducer

Combine Function

Optional Combine function

- Analogous purpose and implementation to the Reduce function
- Objective
 - Decrease the amount of intermediate data ⇒
 i.e. decrease the amount of data transferred to Reducers
- Executed locally by Mapper before the shuffling phase
- Only works for <u>commutative and associative</u> functions!

Combine Function



Intermediate key-value pairs

Counters

- Allow to track the progress of a MapReduce job in real time
 - Predefined counters
 - E.g. numbers of launched / finished Map / Reduce tasks, parsed input key-value pairs, ...
 - Custom counters (user-defined)
 - Can be associated with any action that a Map or Reduce function does

Fault tolerance

When a large number of nodes process a large number of data
 ⇒ fault tolerance is necessary

Worker failure

- Master periodically pings every worker; if no response is received in a certain amount of time, master marks the worker as failed
- All its tasks are reset back to their initial idle state and become eligible for rescheduling on other workers

Master failure

- Strategy A periodic checkpoints are created; if master fails, a new copy can then be started
- Strategy B master failure is considered to be highly unlikely; users simply resubmit unsuccessful jobs

Stragglers

- Straggler = node that takes unusually long time to complete a task it was assigned
- Solution
 - When a MapReduce job is close to completion, the master schedules backup executions of the remaining in-progress tasks
 - A given task is considered to be completed whenever either the primary or the backup execution completes

Task granularity

- Intended numbers of Map and Reduce tasks
- Practical recommendation (by Google)
 - Map tasks
 - Choose the number so that each individual Map task has roughly 16 – 64 MB of input data
 - Reduce tasks
 - Small multiple of the number of worker nodes we expect to use
 - Note also that the output of each Reduce task ends up in a separate output file

Additional Examples

URL access frequency

- Input: HTTP server access logs
- Map: parses a log, emits (accessed URL, 1) pairs
- Reduce: computes and emits the sum of the associated values
- Output: overall number of accesses to a given URL

Inverted index

- Input: text documents containing words
- Map: parses a document, emits (word, document ID) pairs
- Reduce: emits all the associated document IDs sorted
- Output: list of documents containing a given word

Additional Examples

Distributed sort

- Input: records to be sorted according to a specific criterion
- Map: extracts the sorting key, emits (key, record) pairs
- Reduce: emits the associated records unchanged

Reverse web-link graph

- Input: web pages with ... tags
- Map: emits (target URL, current document URL) pairs
- Reduce: emits the associated source URLs unchanged
- Output: list of URLs of web pages targeting a given one

Additional Examples

Sources of links between web pages

```
/**
 * Map function
 * @param key Source web page URL
 * @param value HTML contents of this web page
 */
map(String key, String value) {
  foreach <a> tag t in value: emit(t.href, key);
}
```

```
/**

* Reduce function

* @param key URL of a particular web page

* @param values List of URLs of web pages targeting this one

*/

reduce(String key, Iterator values) {

emit(key, values);
}
```

Use Cases: General Patterns

Counting, summing, aggregation

 When the overall number of occurrences of certain items or a different aggregate function should be calculated

Collating, grouping

 When all items belonging to a certain group should be found, collected together or processed in another way

Filtering, querying, parsing, validation

 When all items satisfying a certain condition should be found, transformed or processed in another way

Sorting

 When items should be processed in a particular order with respect to a certain ordering criterion

Use Cases: Real-World Problems

Just a few real-world examples...

- Risk modeling, customer churn
- Recommendation engine, customer preferences
- Advertisement targeting, trade surveillance
- Fraudulent activity threats, security breaches detection
- Hardware or sensor network failure prediction
- Search quality analysis
- ...



Open-source software framework

- http://hadoop.apache.org/
- Distributed storage and processing of very large data sets on clusters built from commodity hardware
 - Implements a distributed file system
 - Implements a MapReduce programming model
- Derived from the original Google MapReduce and GFS
- Developed by Apache Software Foundation
- Implemented in Java
- Operating system: cross-platform
- Initial release in 2011

Modules

- Hadoop Common
 - Common utilities and support for other modules
- Hadoop Distributed File System (HDFS)
 - High-throughput distributed file system
- Hadoop Yet Another Resource Negotiator (YARN)
 - Cluster resource management
 - Job scheduling framework
- Hadoop MapReduce
 - YARN-based implementation of the MapReduce model

Hadoop-related projects

- Apache Cassandra wide column store
- Apache HBase wide column store
- Apache Hive data warehouse infrastructure
- Apache Avro data serialization system
- Apache Chukwa data collection system
- Apache Mahout machine learning and data mining library
- Apache Pig framework for parallel computation and analysis
- Apache ZooKeeper coordination of distributed applications
- ..

Real-world Hadoop users

- Facebook internal logs, analytics, machine learning, 2 clusters 1100 nodes (8 cores, 12 TB storage), 12 PB 300 nodes (8 cores, 12 TB storage), 3 PB
- LinkedIn 3 clusters 800 nodes (2×4 cores, 24 GB RAM, 6×2 TB SATA), 9 PB 1900 nodes (2×6 cores, 24 GB RAM, 6×2 TB SATA), 22 PB 1400 nodes (2×6 cores, 32 GB RAM, 6×2 TB SATA), 16 PB
- **Spotify** content generation, data aggregation, reporting, analysis 1650 nodes, 43000 cores, 70 TB RAM, 65 PB, 20000 daily jobs
- Yahoo! 40000 nodes with Hadoop, biggest cluster
 4500 nodes (2×4 cores, 16 GB RAM, 4×1 TB storage), 17 PB

HDFS

Hadoop Distributed File System



- Open-source, high quality, cross-platform, pure Java
- Highly scalable, high-throughput, fault-tolerant
- Master-slave architecture
- Optimal applications
 - MapReduce, web crawlers, data warehouses, ...

HDFS: Assumptions

Data characteristics

- Large data sets and files
- Streaming data access
- Batch processing rather than interactive users
- Write-once, read-many

Fault tolerance

- HDFS cluster may consist of thousands of nodes
 - Each component has a non-trivial probability of failure
- ullet \Rightarrow there is always some component that is non-functional
 - I.e. failure is the norm rather than exception, and so
 - automatic failure detection and recovery is essential

HDFS: File System

Logical view: Linux-based hierarchical file system

- Directories and files
- Contents of files is divided into blocks
 - Usually 64 MB, configurable per file level
- User and group permissions
- Standard operations are provided
 - Create, remove, move, rename, copy, ...

Namespace

- Contains names of all directories, files, and other metadata
 - I.e. all data to capture the whole logical view of the file system
- Just a <u>single namespace</u> for the entire cluster

HDFS: Cluster Architecture

Master-slave architecture

- Master: NameNode
 - Manages the file system namespace
 - Manages file blocks (mapping of logical to physical blocks)
 - Provides the user interface for all the operations
 - Create, remove, move, rename, copy, ... file or directory
 - Open and close file
 - Regulates access to files by users
- Slave: DataNode
 - Physically stores file blocks within the underlying file system
 - Serves read/write requests from users
 - I.e. user data never flows through the NameNode
 - Has no knowledge about the file system

HDFS: Replication

Replication = maintaining of multiple copies of each file block

- Increases read throughput, increases fault tolerance
- Replication factor (number of copies)
 - Configurable per file level, usually 3

Replica placement

- Critical to reliability and performance
- Rack-aware strategy
 - Takes the physical location of nodes into account
 - Network bandwidth between the nodes on the same rack is greater than between the nodes in different racks
- Common case (replication factor 3):
 - Two replicas on two different nodes in a local rack
 - Third replica on a node in a different rack

HDFS: NameNode

How the NameNode Works?

- FsImage data structure describing the whole file system
 - Contains: namespace + mapping of blocks + system properties
 - Loaded into the system memory (4 GB RAM is sufficient)
 - Stored in the local file system, periodical checkpoints created
- EditLog transaction log for all the metadata changes
 - E.g. when a new file is created, replication factor is changed, ...
 - Stored in the local file system
- Failures
 - When the NameNode starts up
 - FsImage and EditLog are read from the disk, transactions from EditLog are applied, new version of FsImage is flushed on the disk, EditLog is truncated

HDFS: DataNode

How each **DataNode** Works?

- Stores physical file blocks
 - Each block (replica) is stored as a separate local file
 - Heuristics are used to place these files in local directories
- Periodically sends HeartBeat messages to the NameNode
- Failures
 - When a DataNode fails or in case of a network partition, i.e. when the NameNode does not receive a HeartBeat message within a given time limit
 - The NameNode no longer sends read/write requests to this node, re-replication might be initiated
 - When a DataNode starts up
 - Generates a list of all its blocks and sends a BlockReport message to the NameNode

HDFS: API

Available application interfaces

- Java API
 - Python access or C wrapper also available
- HTTP interface
 - Browsing the namespace and downloading the contents of files
- FS Shell command line interface
 - Intended for the user interaction
 - Bash-inspired commands
 - E.g.:
 - hadoop fs -ls /
 - hadoop fs -mkdir /mydir

Hadoop MapReduce

Hadoop MapReduce



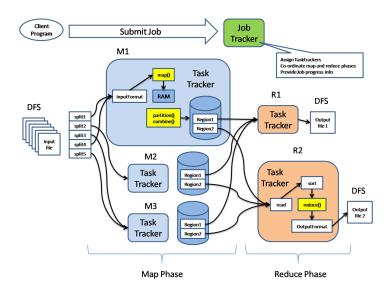
- MapReduce programming model implementation
- Requirements
 - HDFS
 - Input and output files for MapReduce jobs
 - YARN
 - Underlying distribution, coordination, monitoring and gathering of the results

Cluster Architecture

Master-slave architecture

- Master: JobTracker
 - Provides the user interface for MapReduce jobs
 - Fetches input file data locations from the NameNode
 - Manages the entire execution of jobs
 - Provides the progress information
 - Schedules individual tasks to idle TaskTrackers
 - Map, Reduce, ... tasks
 - Nodes close to the data are preferred
 - Failed tasks or stragglers can be rescheduled
- Slave: TaskTracker
 - Accepts tasks from the JobTracker
 - Spawns a separate JVM for each task execution
 - Indicates the available task slots via HearBeat messages

Execution Schema



Java Interface

Mapper class

- Implementation of the map function
- Template parameters
 - KEYIN, VALUEIN types of input key-value pairs
 - KEYOUT, VALUEOUT types of intermediate key-value pairs
- Intermediate pairs are emitted via context.write(k, v)

Java Interface

Reducer class

- Implementation of the reduce function
- Template parameters
 - KEYIN, VALUEIN types of intermediate key-value pairs
 - KEYOUT, VALUEOUT types of output key-value pairs
- Output pairs are emitted via context.write(k, v)

```
class MyReducer extends Reducer<KEYIN, VALUEIN, KEYOUT, VALUEOUT> {
   @Override
   public void reduce(KEYIN key, Iterable<VALUEIN> values, Context context)
      throws IOException, InterruptedException
   {
        // Implementation
   }
}
```

Example

Word Frequency

- Input: Documents with words
 - Files located at /home/input HDFS directory
- Map: parses a document, emits (word, 1) pairs
- Reduce: computes and emits the sum of the associated values
- Output: overall number of occurrences for each word
 - Output will be written to /home/output

MapReduce job execution

hadoop jar wc.jar WordCount /home/input /home/output

Example: Mapper Class

```
public class WordCount {
 public static class MyMapper
   extends Mapper<Object, Text, Text, IntWritable>
    private final static IntWritable one = new IntWritable(1):
   private Text word = new Text();
   Onverride
   public void map(Object key, Text value, Context context)
     throws IOException, InterruptedException
     StringTokenizer itr = new StringTokenizer(value.toString());
     while (itr.hasMoreTokens()) {
        word.set(itr.nextToken()):
        context.write(word, one);
```

Example: Reducer Class

```
public class WordCount {
 public static class MyReducer
   extends Reducer < Text. IntWritable. Text. IntWritable>
    private IntWritable result = new IntWritable():
   @Override
   public void reduce(Text key, Iterable<IntWritable> values,
     Context context) throws IOException, InterruptedException
     int sum = 0:
     for (IntWritable val : values) {
        sum += val.get();
     result.set(sum):
     context.write(key, result);
```

Lecture Conclusion

MapReduce criticism

- MapReduce is a step backwards
 - Does not use database schema
 - Does not use index structures
 - Does not support advanced query languages
 - Does not support transactions, integrity constraints, views, ...
 - Does not support data mining, business intelligence, ...
- MapReduce is not novel
 - Ideas more than 20 years old and overcome
 - Message Passing Interface (MPI), Reduce-Scatter

The end of MapReduce?